

New Ideas Track: Testing MapReduce-Style Programs

Christoph Csallner, Leonidas Fegaras, Chengkai Li
Computer Science and Engineering Department
University of Texas at Arlington (UTA)

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Since 2004: Many MapReduce systems, papers & users

- Google MapReduce [OSDI 2004] > 2,000 cit.
- Apache/Yahoo! Hadoop
 - <http://wiki.apache.org/hadoop/PoweredBy>
- Microsoft Dryad [EuroSys 2007] > 500 cit.
 - <http://research.microsoft.com/en-us/projects/dryad/>
- Apache/Yahoo! Pig [SIGMOD 2008] > 400 cit.
 - <https://cwiki.apache.org/confluence/display/PIG/PoweredBy>
- Apache/Facebook Hive [VLDB 2009]
 - <https://cwiki.apache.org/confluence/display/Hive/PoweredBy>

MapReduce programming model

- **Programmer implements sequential code**
 - Two functions: map and reduce
 - For example, in sequential Java code
- **System distributes, schedules, handles faults**
 - Invokes map **on many nodes in parallel**
 - Collects and re-distributes intermediate results
 - Invokes reduce **on many nodes in parallel**
- **Programmer can focus on problem domain**

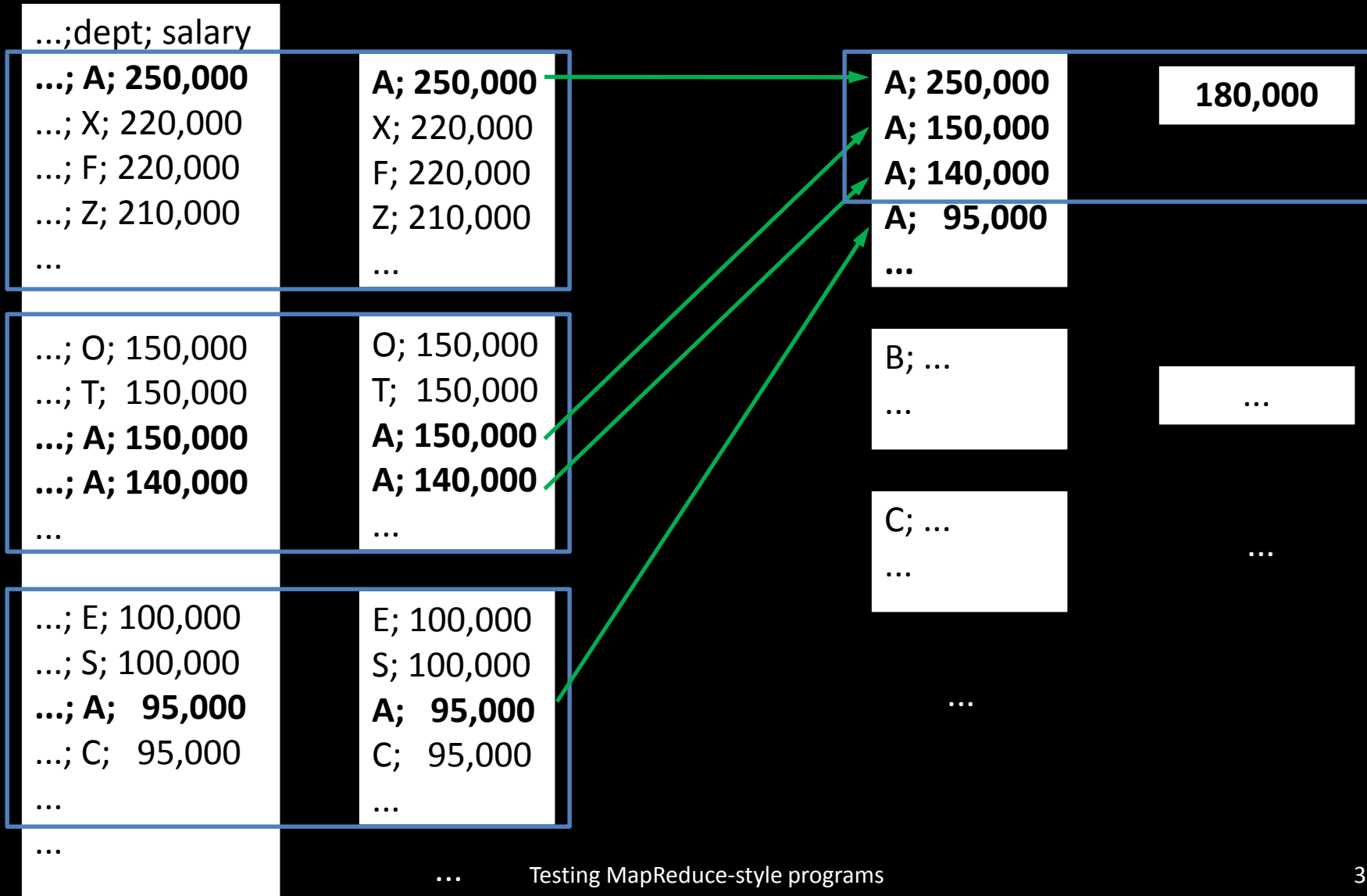
Input

Map:
(key;value)*

Group
By Key

Reduce:
avg of first 3

Output



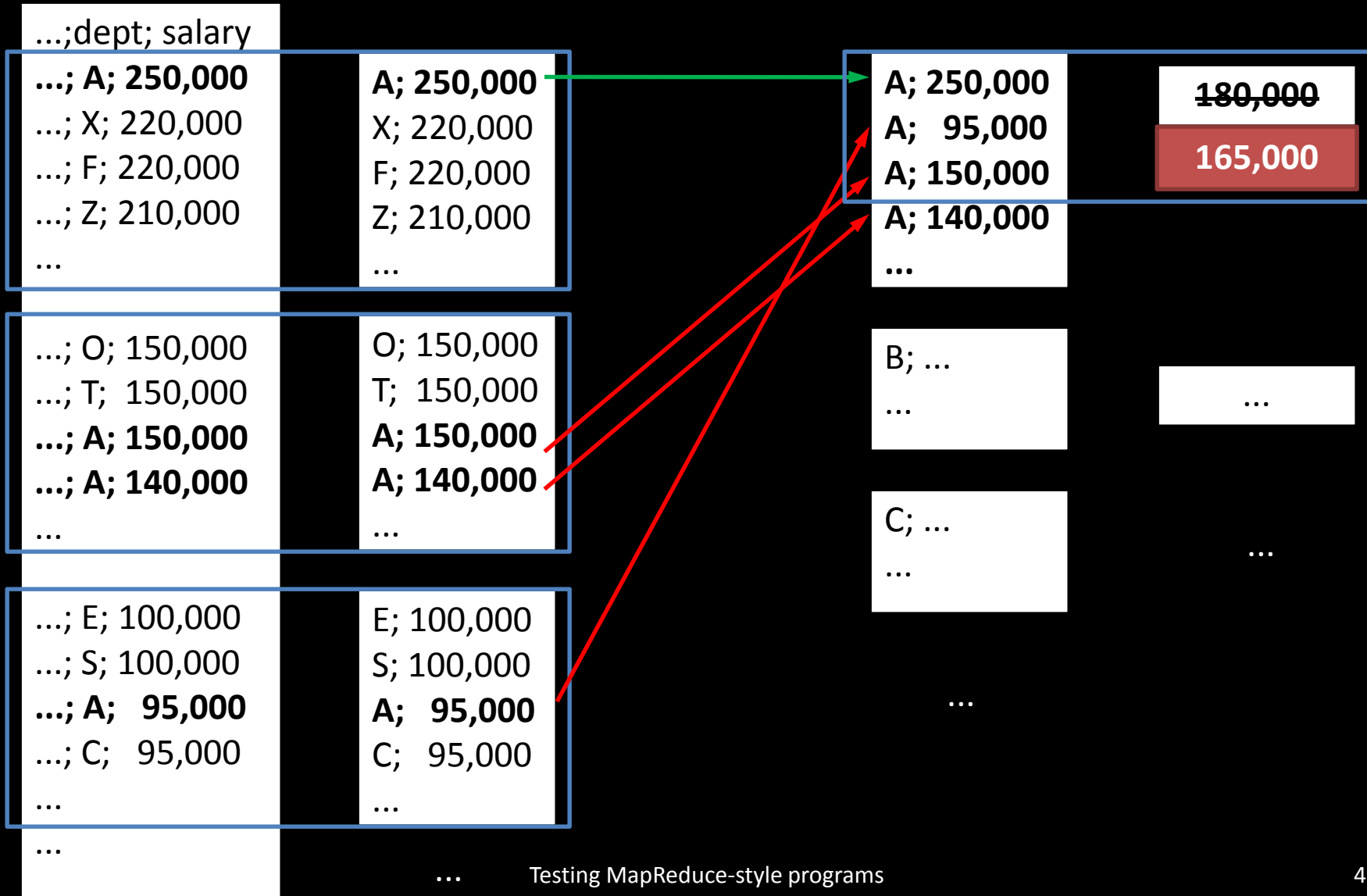
Input

Map:
(key;value)*

Group
By Key

Reduce:
avg of first 3

Output



Example

bug:

```
/* Report avg of top-3 salaries, if avg>100k */  
public void reduce(String dept, Iterator<Integer> salaries) {  
    int sum = 0; int i = 0;  
    while (salaries.hasNext() && i<3) {  
        sum += salaries.next();  
        i += 1;  
    }  
    emit( (i>0 && sum/i > 100000)? sum/i : -1);  
}
```

- Code depends on order of salaries, just uses first-3
- Programmer may be confused by order of salaries in input files, that order is not maintained
- Bug, possibly because MapReduce systems have built-in ordering, but not always use them

User reduce program has to satisfy correctness conditions

- Reduce must not rely on a particular order:
- For each input list of values L,
for each permutation P:
 $\text{reduce}(\text{key}, L) == \text{reduce}(\text{key}, P(L))$
- Program also has to satisfy other MapReduce-specific correctness conditions
- Current tools do not check these conditions

Goal: Find such bugs automatically

- Find an input list of values L and a permutation P :
 $\text{reduce}(\text{key}, L) \neq \text{reduce}(\text{key}, P(L))$
- Current tools do not find such bugs
- There are many input lists and permutations
 - Trying all of them is impossible

Example

bug:

```
/* Report avg of top-3 salaries, if avg>100k */  
public void reduce(String dept, Iterator<Integer> salaries) {  
    int sum = 0; int i = 0;  
    while (salaries.hasNext() && i<3) {  
        sum += salaries.next();  
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    }  
    emit( (i>0 && sum/i > 100000)? sum/i : -1);  
}
```

- Need specific list of salaries & permutation
 - List of more than 3 elements
 - Average of first 3 elements > 100k
 - Permutation has to swap element at position ≤ 3 with element at position > 3

Observations

- Example MapReduce programs are typically small and contain few execution paths
 - How do industrial MapReduce programs look like?
- Dynamic symbolic execution may be a good fit
 - Heavy-weight but precise analysis
 - Systematically explores all execution paths
 - Well-suited for reasoning about few paths
- `reduce(key, L)`, `reduce(key, P(L))` may trigger different execution paths
 - Not enough to analyze one path at a time

Check correctness conditions with dynamic symbolic execution

1. Derive symbolic path condition, return value
2. Maintain them in an indexed execution tree
 - Index leaf nodes by length of input list
 - Sibling(path): Triggered by input list of same length
3. Encode potential violation of correctness condition in constraint system
 - Solving constraints with off-the-shelf constraint solver yields concrete input values L and permutation P
4. Convert solution to test case, run, confirm violation

Encode correctness conditions in symbolic program constraints

```
// Permutation P as a function: 0 → p[0], 1 → p[1], ..  
// Symbolic list L = L[0], L[1], ..    P(L) = L[p[0]], L[p[1]], ..  
SymbolicInt[] p ← SymbolicIndices; // distinct list positions  
Assert PathCond; // e.g.: L[0]==5  
Assert SubstituteIndices(SiblingPath, p); // e.g.: L[p[0]]==5  
  
// Find a concrete list + a concrete permutation such that:  
// reduce(key, list) ≠ reduce(key, permutation(list))  
Assert Result ≠ SubstituteIndices(SiblingResult, p);
```

Input length heuristic

- Pick “representative” input lengths
- Initially: $|L| := 2$
 - For shorter lists: $L == P(L)$
- Binary back-off scheme
 - Each subsequent iteration doubles length of L

Conclusions

- New programming paradigm with new bugs
 - To produce deterministic results, a MapReduce system requires user programs to satisfy certain high-level correctness conditions
 - Neither MapReduce execution systems nor tools check these conditions
- Proposed approach:
 - Encode MapReduce correctness conditions in symbolic program constraints
 - Check correctness conditions at runtime

References

- **[OSDI 2004]** J. Dean and S. Ghemawat. *MapReduce: Simplified data processing on large clusters*. In Proc. 6th USENIX Symposium on Operating Systems Design and Implementation, pages 137—150.
- **[EuroSys 2007]** M. Isard, M. Budiu, Y. Yu, A. Birrell, and D. Fetterly. *Dryad: Distributed data-parallel programs from sequential building blocks*. In Proc. 2nd ACM SIGOPS European Conference on Computer Systems, pages 59—72.
- **[SIGMOD 2008]** C. Olston, B. Reed, U. Srivastava, R. Kumar, and A. Tomkins. *Pig latin: A not-so-foreign language for data processing*. In Proc. 34th ACM SIGMOD International Conference on Management of Data, pages 1099—1110.
- **[CACM2008]** J. Dean and S. Ghemawat. MapReduce: Simplified data processing on large clusters. *Communications of the ACM*, 51(1):107—113.
- **[VLDB 2009]** A. Thusoo, J. S. Sarma, N. Jain, Z. Shao, P. Chakka, S. Anthony, H. Liu, P. Wycko, and R. Murthy. *Hive: A warehousing solution over a map-reduce framework*. *Proc. VLDB Endowment*, 2(2):1626—1629.

Questions

MapReduce used for variety of jobs

- Process “web-scale” data (PB = peta-byte = 10^{15})
 - Run on many machines in parallel
- Google: Process 20 PB per day [CACM2008]
 - 10k programs build search index, process text, graphs, etc.
- New York Times: Convert 4TB of articles to PDF
 - <http://open.blogs.nytimes.com/2007/11/01/self-service-prorated-super-computing-fun/>
- Yahoo!: Sort TB in 209 seconds: <http://sortbenchmark.org/>
 - “First time that either a Java or an open source program has won this challenge” [<http://hadoop.apache.org/>]
- Facebook: Hive-based data warehouse

MapReduce \neq map-reduce

- MapReduce:
 - Inspired by functional programming map-reduce
 - But different 😊
- For detailed comparison, see:
 - Ralf Lämmel. *Google's MapReduce programming model — Revisited*. *Science of Computer Programming* 68(3): 208—237. Oct. 2007.

MapReduce correctness condition 2: Optional combine function

- Combine: programmer-defined sequential code
 - Similar to map and reduce
- May be invoked on Map node, after map
 - Locally “pre-reduce” results, by key
 - Reduce transmission overhead to “real reduce”
- System can invoke combine 0—n times
 - Must not affect semantics
- Similar approach:
 - Encode in symbolic path condition, result value