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# Design and Analysis of Algorithms

CSE 5311

Lecture 1 Administration & Introduction

Junzhou Huang, Ph.D.

Department of Computer Science and Engineering

# Administration

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- **Course CSE5311**

- What: Design and Analysis of Algorithms
- When: Friday 1:00 ~ 3:50pm
- Where: ERB 130
- Who: Junzhou Huang (Office ERB 650) [jzhuang@uta.edu](mailto:jzhuang@uta.edu)
- Office Hour: Friday 3:50 ~ 5:50pm and/or appointments
- Homepage: <http://ranger.uta.edu/~huang/teaching/CSE5311.htm>  
(**You're required to check this page regularly**)

- **Lecturer**

- PhD in CS from Rutgers, the State University of New Jersey
- Research areas: machine learning, computer vision, medical image analysis and bioinformatics

- **GTA**

- Saiyang Na (Office ERB 403), [sxn3892@mavs.uta.edu](mailto:sxn3892@mavs.uta.edu)
- Office hours: Friday 10:00am ~ 12:00pm and/or appointments

# Study Materials

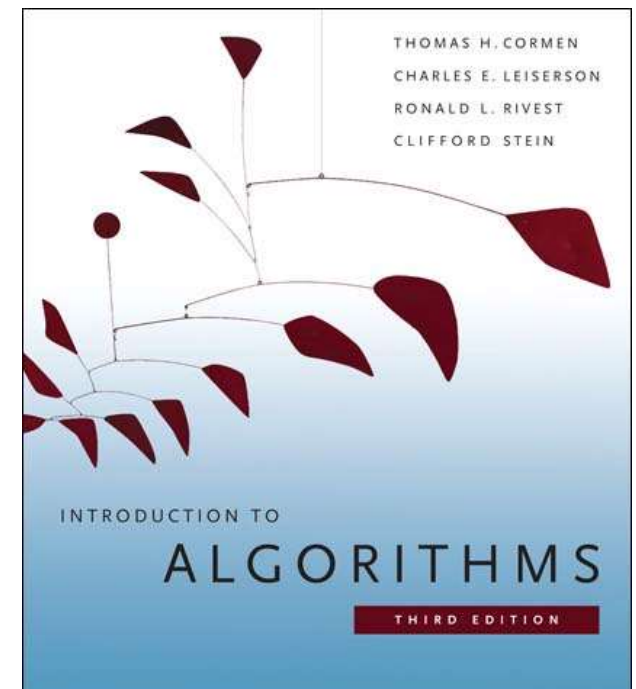
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- **Prerequisites**

- Algorithms and Data Structure (CSE 2320)
- Theoretical Computer Science (CSE 3315)
- What this really means:
  - You have working experience s on software development.
  - You know compilation process and programming
  - Elementary knowledge of math and algorithms

- **Text book**

- [Thomas H. Cormen](#), [Charles E. Leiserson](#), [Ronald L. Rivest](#) and [Clifford Stein](#), **Introduction to Algorithms**, third edition
- <https://mitpress.mit.edu/books/introduction-algorithms>



# Study Materials

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- **Text book**
  - We will not cover all the chapters of the book
  - We will not cover all sections of the covered chapters
  - We will not fully follow the order of the book
  - The contents uncovered in slides/lectures are optional
- **Recommended Reference**
  - Robert Sedgwick, **Algorithms in C**. Addison - Wesley, 1990
- **Acknowledgments**
  - Class notes partially based on 5311 classes taught at UTA in prior years
  - Material from textbook site
  - Lots of material available on the web (via google search, wikipedia)

# Grading

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- **Distribution**

- 20%            Projects
  - 35%            Midterm Exam
  - 40%            Final Exam
  - 5%             Class Participation
- 
- 100%

- **Attention**

- Homework is as important as any other aspects of your grade!
- Attendance though not mandatory, but is HIGHLY encouraged.
- The university makeup policy will be strictly adhered to. Generally, no make-up exams/quizzes except for university sanctioned reasons.
- When missing an exam/quiz due to unavoidable circumstances, PLEASE notify the instructor and **request a makeup approval ahead of time.**

# Final Grade

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- **Final Letter Grade**

- [90 100] --- A
- [80 90) --- B
- [70 80) --- C
- [60 70) --- D
- [00 60) --- F

- **Attention**

- Final letter grades will be assigned based on absolute percentage
- [ ] denotes inclusion and ( ) denotes exclusion.
- The instructor reserves the right to move the thresholds down based on the distribution of final percentages, but they will not move up.

# Assignments

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- **Homework assignments**
  - Assigned in class, typically due one week later at the start of next lecture
  - Solutions will be posted soon after the due for reference
- **Projects**
  - Implementing and comparing the learnt algorithms in the class.
  - Each group has two members at most.
  - Each group should submit a project report and do a presentation in the class
  - Each member in the group should mention her/his contribution respectively
- **Collaboration**
  - You may discuss assignments with others, but must write up them individually. Please identify collaborators on your assignment cover sheet
  - Failure to comply with this policy is a violation of academic integrity
- **Start early! Start early! Start early !!!**

# Information

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- **Course Webpage**
  - Check the web page regularly (2 times per week).
  - Announcements, assignments, and lecture notes will be posted there.
- **Grade Appeal**
  - You may appeal the grade in writing (email) within 5 class days.
  - Appealed to the appropriate GTA firstly, then to the instructor, if necessary.
  - Please refer to the UTA Catalog for the detailed guide of grade appeals.
- **Drop Policy**
  - The university withdrawal policy will be strictly adhered to.
- **Others**
  - Accommodating students with disabilities
  - Student Support Services
  - Etc.



# Questions

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# Course Overview

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- **What is it?**
  - Algorithms
  - Design and Analysis
- **Why is a CS course?!?**
  - The key of the computer science
- **Will I really ever use this stuff again?**
  - Definitely, analysis and design
  - Necessary knowledge for a CS student
  - You may not become a professional algorithm designer but you definitely need know how to analysis and design the algorithms for the problems in your future career and even in your life
- **How to succeed in this course?**

# Why Are You In This Class?

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- **Something interesting about you**
  - Why you picked your major?
  - Life Plan
- **To learn background in order to take more advanced classes in computer science and engineering**
  - Database, big data analytics, compiler, Computer Network, Embedded Systems, artificial intelligent, machine learning, data mining, computer vision, etc.
- **Understand the effect of an algorithm on the code you write or read**
  - Learn how to efficiently use and control the computer
- **To have the necessary background to understand innovations in intelligent design or related others**
  - Your desktop, laptop, ipad, iphone, google search, facebook social network, etc.
- **Necessary to become a professional algorithm designer?**
  - But to be able to read and understand
  - To be able to understand innovative ideas

# What?

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- *The theoretical study of design and analysis of computer algorithms*
- **Basic goals for an algorithm**
  - Always correct
  - Always terminates
- **Our class: performance**
  - Performance often draws the line between what is possible and what is impossible.
- **Design and Analysis of Algorithms**
  - **Analysis:** predict the cost of an algorithm in terms of resources and performance
  - **Design:** design algorithms which minimize the cost

# Machine Model

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- **Generic Random Access Machine (RAM)**
  - Executes operations sequentially
  - Set of primitive operations: Arithmetic, Logical, Comparisons, Function calls
- **Simplifying assumption**
  - All operations cost 1 unit
  - Eliminates dependence on the speed of our computer
  - Otherwise impossible to verify and to compare

# The Problem of Sorting

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*Input:* sequence  $\langle a_1, a_2, \dots, a_n \rangle$  of numbers.

*Output:* permutation  $\langle a'_1, a'_2, \dots, a'_n \rangle$  such that  $a'_1 \leq a'_2 \leq \dots \leq a'_n$ .

**Example:**

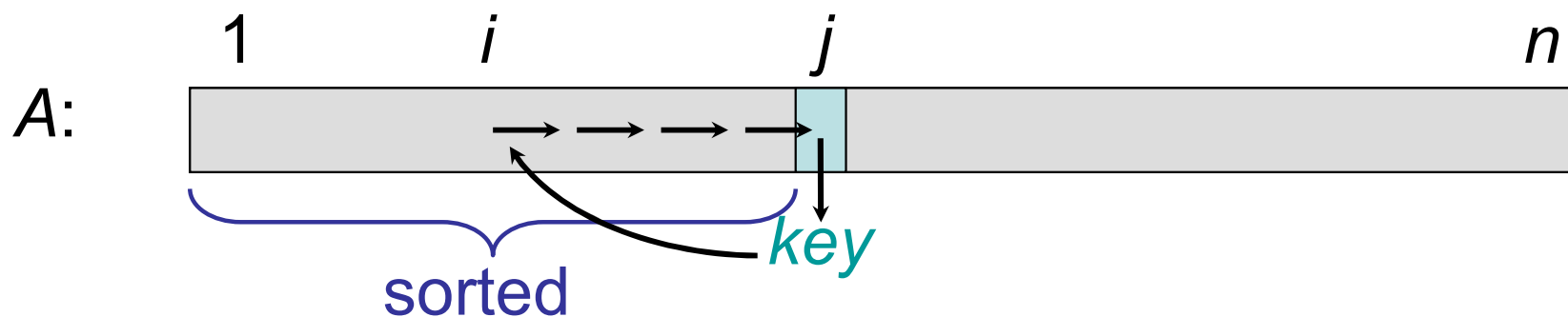
*Input:* 8 2 4 9 3 6

*Output:* 2 3 4 6 8 9

# Insertion sort

“pseudocode”

```
INSERTION-SORT ( $A, n$ )    ▷  $A[1 \dots n]$   
  for  $j \leftarrow 2$  to  $n$   
    do  $key \leftarrow A[j]$   
       $i \leftarrow j - 1$   
      while  $i > 0$  and  $A[i] > key$   
        do  $A[i+1] \leftarrow A[i]$   
           $i \leftarrow i - 1$   
       $A[i+1] = key$ 
```



# Example of Insertion Sort

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8   2   4   9   3   6

The image shows a horizontal sequence of six numbers: 8, 2, 4, 9, 3, and 6. The number 2 is enclosed in a bright yellow circle, indicating it is the current element being processed in the insertion sort algorithm. The other numbers are in a teal color.



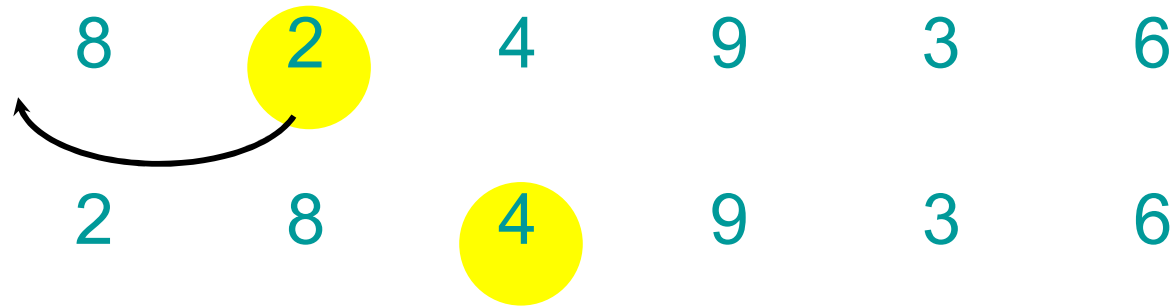
# Example of Insertion Sort

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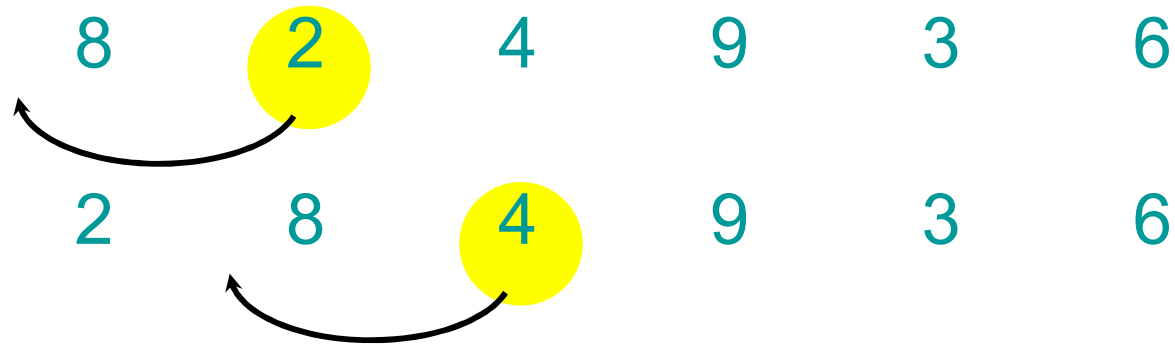
# Example of Insertion Sort

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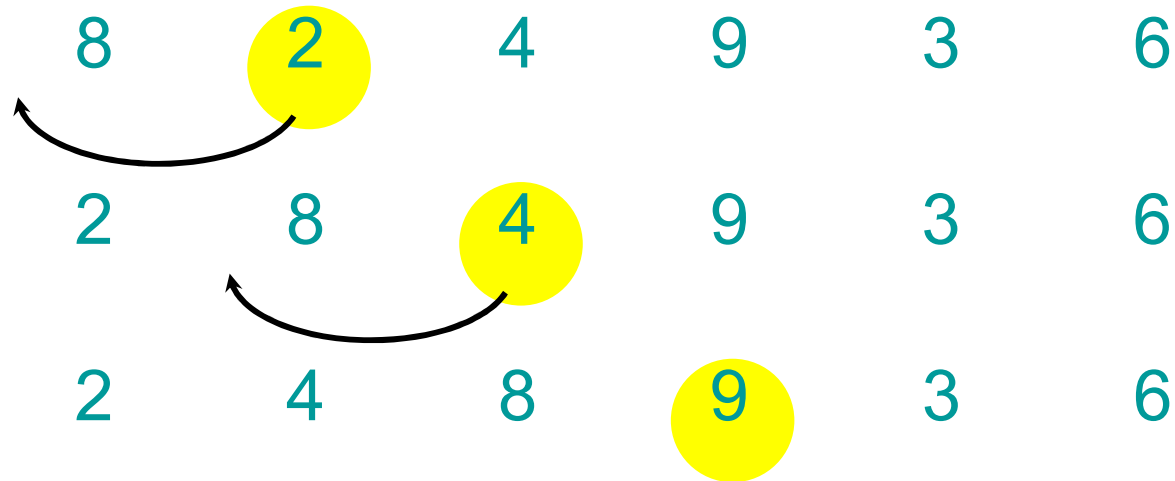
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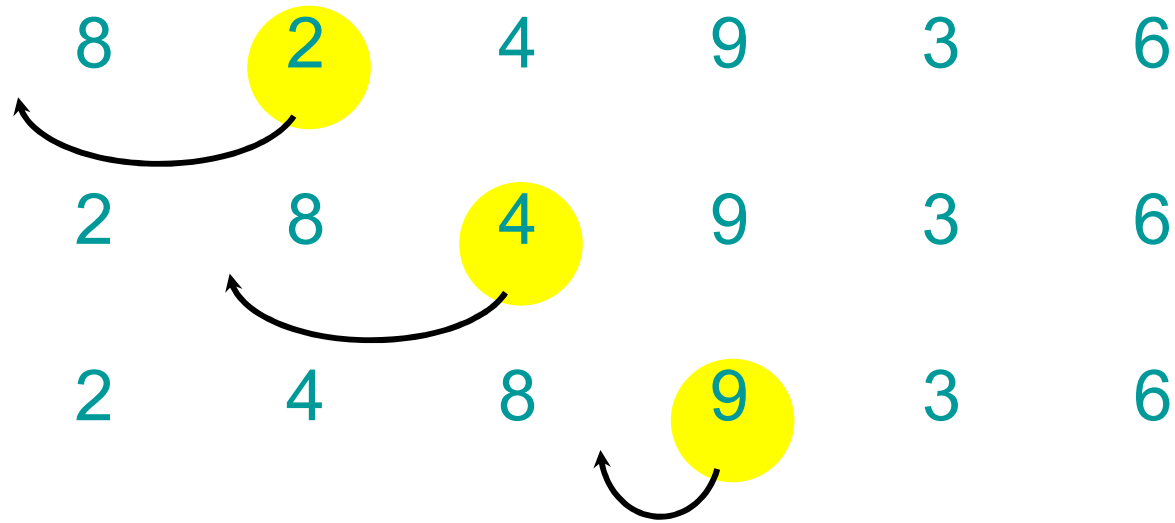
# Example of Insertion Sort

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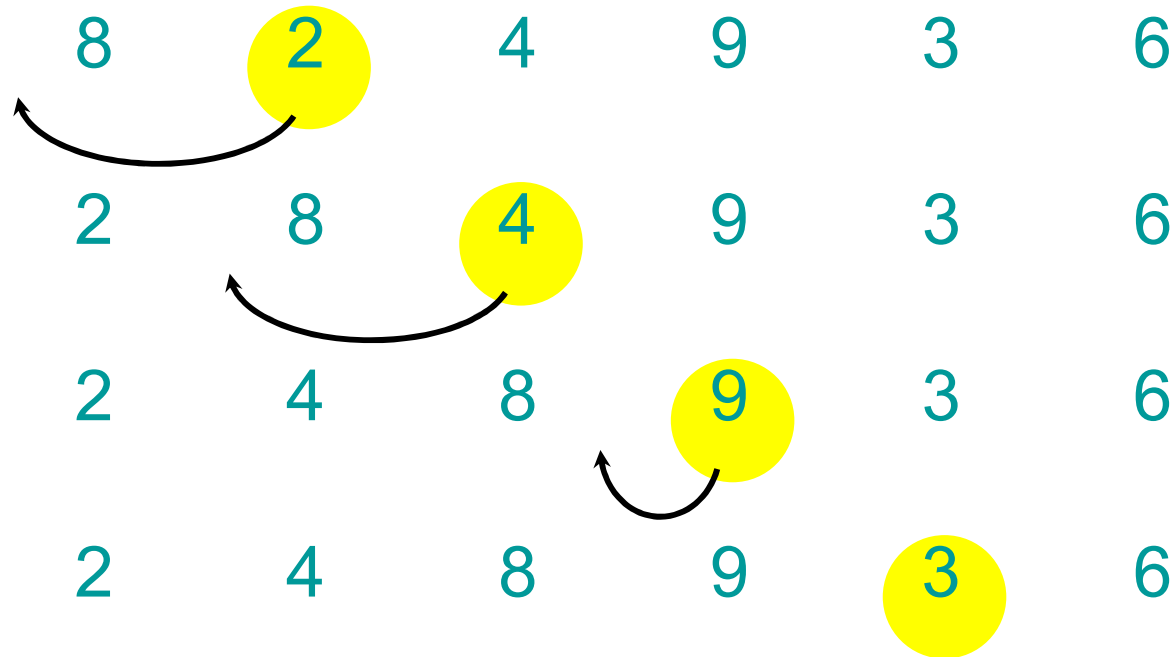
# Example of Insertion Sort

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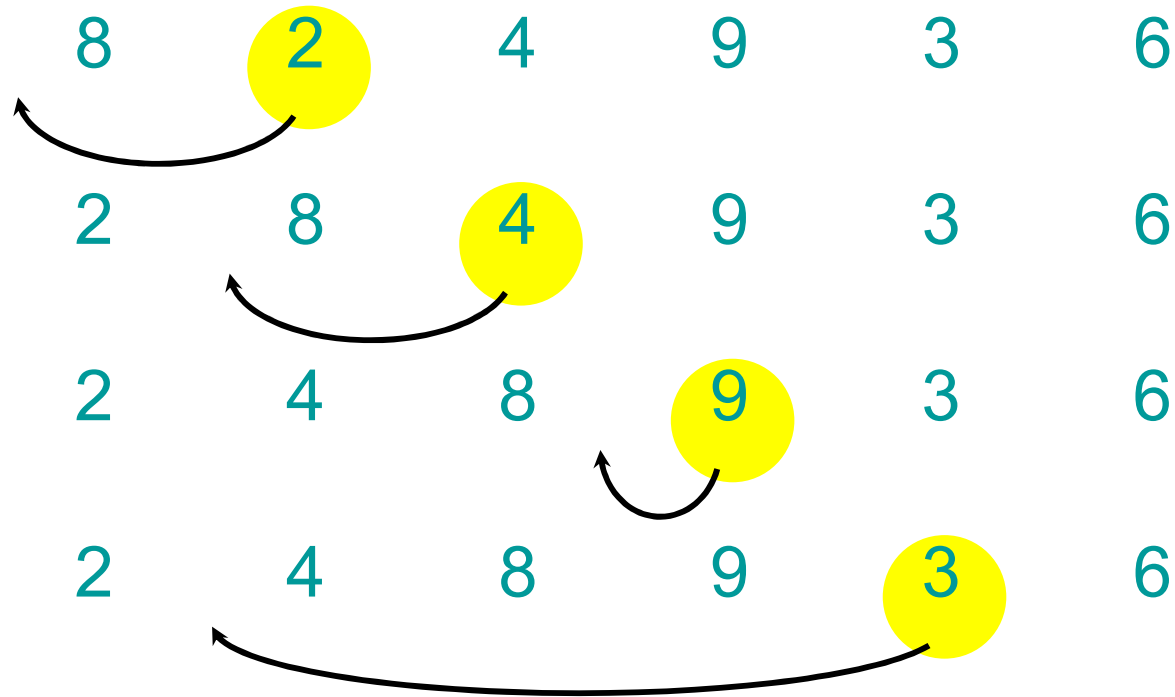
# Example of Insertion Sort

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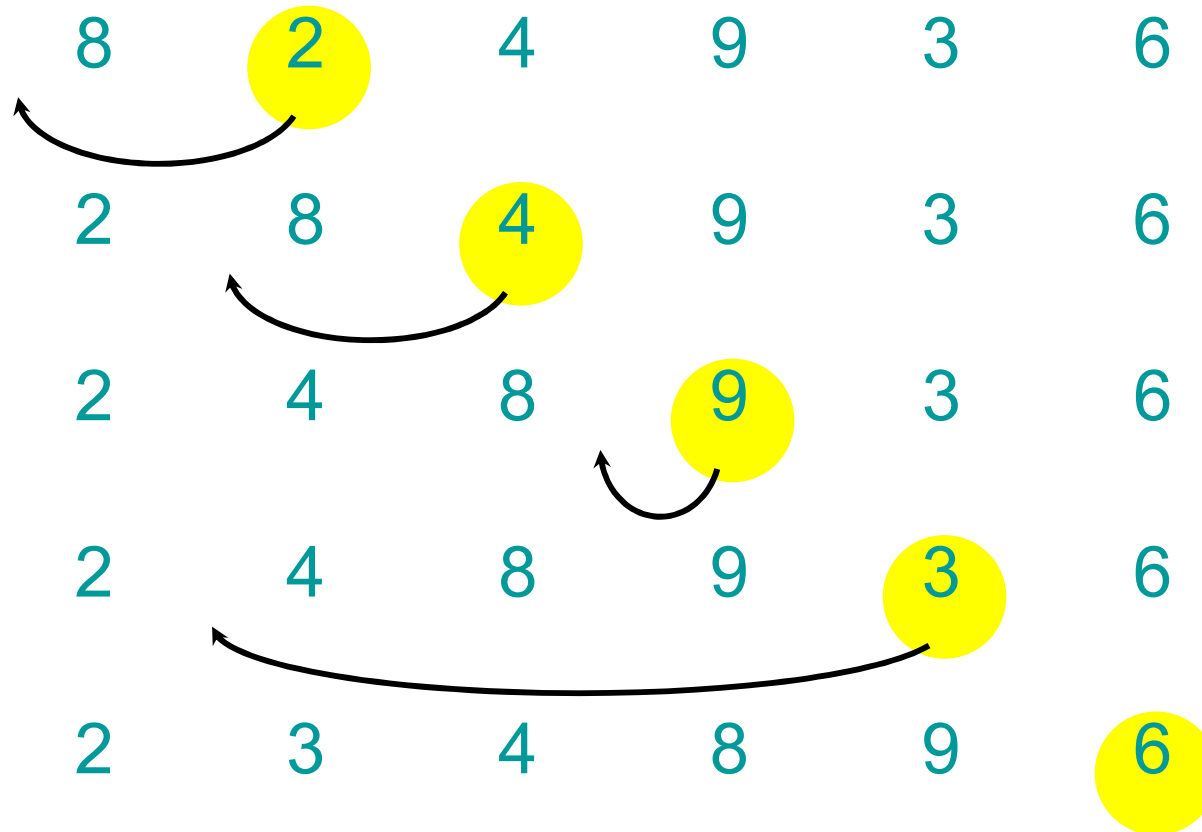
# Example of Insertion Sort

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# Example of Insertion Sort

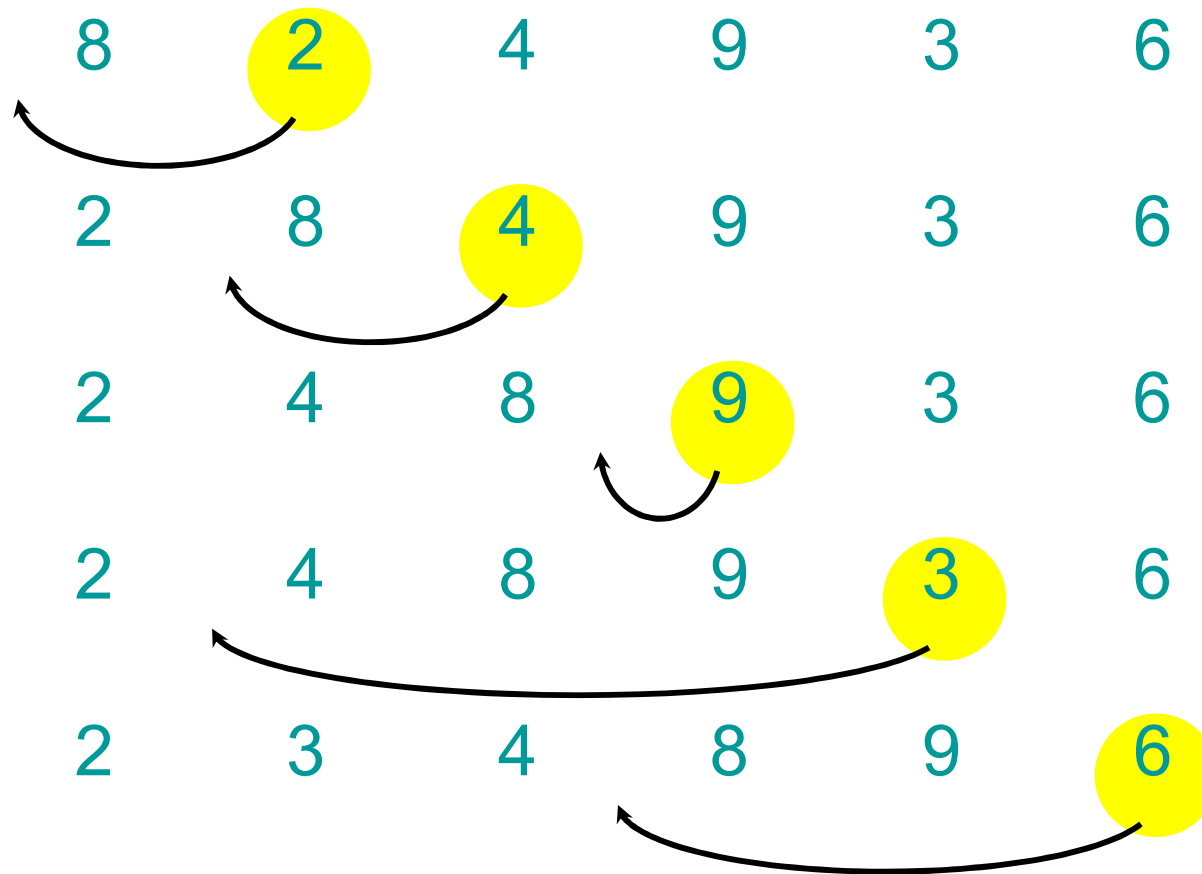
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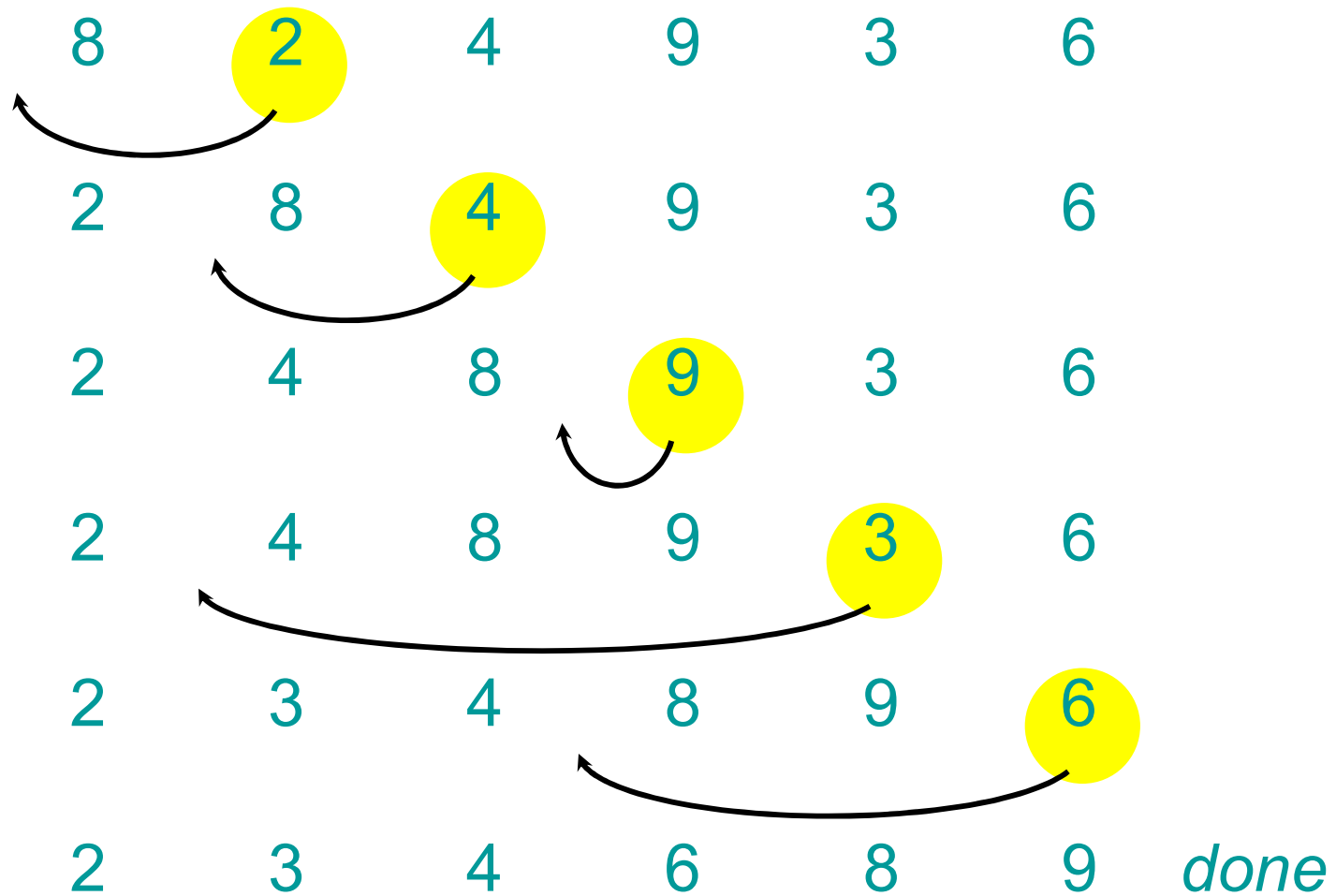
# Example of Insertion Sort

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# Example of Insertion Sort

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# Running Time

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- **Running Time**
  - Depends on the input
  - An already sorted sequence is easier to sort.
- **Major Simplifying Convention**
  - Parameterize the running time by the size of the input, since short sequences are easier to sort than long ones.
  - $T_A(n)$  = time of  $A$  on length  $n$  inputs. Generally, we seek upper bounds on the running time, to have a guarantee of performance.
- **Kinds of analyses**
  - **Worst-case:** (usually)  $T(n)$  = maximum time of algorithm on any input of size  $n$
  - **Average-case:** (sometimes)  $T(n)$  = expected time of algorithm over all inputs of size  $n$ . Need assumption of statistical distribution of inputs.
  - **Best-case:** (Never) Cheat with a slow algorithm that works fast on some input.

# Measuring Algorithm Complexity

- How long does it take to execute a program?
  - Efficient algorithm, the better
- How long does it take to go from point A to Point B

- **Need to know:**
  - Speed: Walk/driving
  - Distance: variable
  - Traffic condition
  - Walk?
  - Driving?



# Analysis

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- **Simplifications**

- Ignore actual and abstract statement costs
- Order of growth is the interesting measure:
  - Highest-order term is what count
  - Doing asymptotic analysis
  - As the input size grows larger it is the high order term that dominates

- **Teaching Goals**

- Show that by knowing more about the underlying algorithm design and analysis, one can be more effective as a computer scientist or engineer.

# Upper Bound Notation

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- **Definition**

- In general, a function  $f(n)$  is  $O(g(n))$  if there exist positive constants  $c$  and  $n_0$  such that  $f(n) \leq c \cdot g(n)$  for all  $n \geq n_0$
- Formally,  $O(g(n)) = \{ f(n): \exists \text{ positive constants } c \text{ and } n_0 \text{ such that } f(n) \leq c \cdot g(n) \forall n \geq n_0$

- **A polynomial of degree  $k$  is  $O(n^k)$**

- **Proof:**

Suppose  $f(n) = b_k n^k + b_{k-1} n^{k-1} + \dots + b_1 n + b_0$

Let  $a_i = |b_i|$

$f(n) \leq a_k n^k + a_{k-1} n^{k-1} + \dots + a_1 n + a_0$

$$\leq n^k \sum a_i \frac{n^i}{n^k} \leq n^k \sum a_i \leq cn^k$$

# Upper Bound Notation

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- **We say InsertionSort's run time is  $O(n^2)$** 
  - Properly we should say run time is *in*  $O(n^2)$ ; Read O as “Big-O”
- **Proof**
  - Suppose runtime is  $an^2 + bn + c$ ;
  - If any of  $a$ ,  $b$ , and  $c$  are less than 0 replace the constant with its absolute value
  - $an^2 + bn + c \leq (a + b + c)n^2 + (a + b + c)n + (a + b + c)$   
 $\leq 3(a + b + c)n^2$   
for  $n \geq 1$ ; Let  $c' = 3(a + b + c)$  and let  $n_0 = 1$
- **Questions**
  - Is InsertionSort  $O(n^3)$ ?
  - Is InsertionSort  $O(n)$ ?

# Lower Bound Notation

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- **Definition**

- In general a function  $f(n)$  is  $\Omega(g(n))$  if  $\exists$  positive constants  $c$  and  $n_0$  such that  $0 \leq c \cdot g(n) \leq f(n) \quad \forall n \geq n_0$

- **We say InsertionSort's run time is  $\Omega(n)$**

- **Proof**

- Suppose run time is  $a \cdot n + b$
- Assume  $a$  and  $b$  are positive (what if  $b$  is negative?)  
 $a \cdot n \leq a \cdot n + b$



# Asymptotic Tight Bound

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- **Asymptotic Tight Bound  $\Theta$**

- A function  $f(n)$  is  $\Theta(g(n))$  if  $\exists$  positive constants  $c_1, c_2$ , and  $n_0$  such that
$$c_1 g(n) \leq f(n) \leq c_2 g(n), \quad \forall n \geq n_0$$
- **Theorem:**  $f(n)$  is  $\Theta(g(n))$  iff  $f(n)$  is both  $O(g(n))$  and  $\Omega(g(n))$
- **Proof:** self practice

- **Other Asymptotic Tight Bounds**

- A function  $f(n)$  is  $o(g(n))$  if  $\exists$  positive constants  $c$  and  $n_0$  such that  $f(n) < c g(n)$   
 $\forall n \geq n_0$
- A function  $f(n)$  is  $\omega(g(n))$  if  $\exists$  positive constants  $c$  and  $n_0$  such that  $c g(n) < f(n)$   
 $\forall n \geq n_0$
- Intuitively,

- |                        |                             |                          |
|------------------------|-----------------------------|--------------------------|
| ▪ $o()$ is like $<$    | ▪ $\omega()$ is like $>$    | ▪ $\Theta()$ is like $=$ |
| ▪ $O()$ is like $\leq$ | ▪ $\Omega()$ is like $\geq$ |                          |

# Course Goals

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- **Teaching Style**

- Algorithm courses have been traditionally taught by following textbooks and covering basic concepts and algorithms for different problems.
- We're going to follow this style and additionally cover some fresh problems and algorithms appeared in FLG interviewing.

- **Teaching Goals**

- Show that by knowing more about the underlying algorithm design and analysis, one can be more effective as a computer scientist or engineer.
- Write programs that are more reliable and efficient for the specific goal.
- Understand how program performance depends on underlying factors including the designed algorithm.
- Learn how to implement an effective and efficient software system according to the request and the available resources (memory, communication, disk, GPU, etc.)

# Course Expectation

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- **What to expect from the course:**
  - Will cover key issues and concepts in class.
  - Recitations will provide review and teach you the ideas you need.
  - Programming Projects (Don't freak out... yet)
  - A mid-term exam and a final exam
  - Practice homework sets
- **What do I expect of you:**
  - Come to class
  - Read the textbook (Listening to me is not good enough)
  - Work through the problems in the textbook (not really homework... but it helps)
  - Do the projects
  - Ask questions (**IMPORTANT**)

# Questions

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