RankSQL:

Query Algebra and Optimization for Relational Top-k Queries

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joint work with

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Ranking (Top-k) Queries

Ranking is an important functionality in many real-world database applications:

- E-Commerce, Web Sources
 Find the best hotel deals by price, distance, etc.
- Multimedia Databases
 Find the most similar images by color, shape, texture, etc.
- Text Retrieval, Search Engine
 Find the most relevant records/documents/pages.
- OLAP, Decision Support
 Find the top profitable customers to send ads.

Example: Trip Planning

Suggest a hotel to stay and a museum to visit:

Select *

From

Hotel h, Museum m

Where

hotel	museum	cheap	close	related	score
h1	m2	0.9	0.7	0.8	2.4
h2	m1	0.6	0.8	0.9	2.3
h1	m3	0.9	0.7	0.6	2.2

h.star=3 AND h.area=m.area

Order By

cheap(h.price) +
close(h.addr, "BWI airport") +
related(m.collection, "dinosaur")

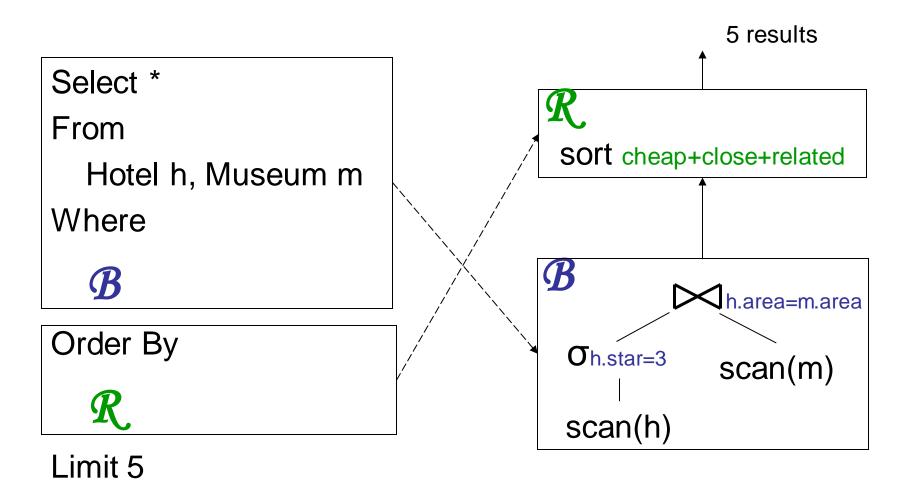


membership dimension: Boolean predicates, Boolean function

order dimension: ranking predicates, monotonic scoring function

Limit 5

Processing Ranking Queries in Traditional RDBMS

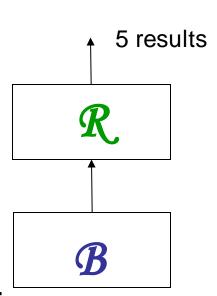


Problems of Traditional Approach

- Naïve Materialize-then-Sort scheme
- Overkill:

total order of all results; only 5 top results are requested.

- Very inefficient:
 - Scan large base tables;
 - Join large intermediate results;
 - Evaluate every ranking on every tuple;
 - Full sorting.

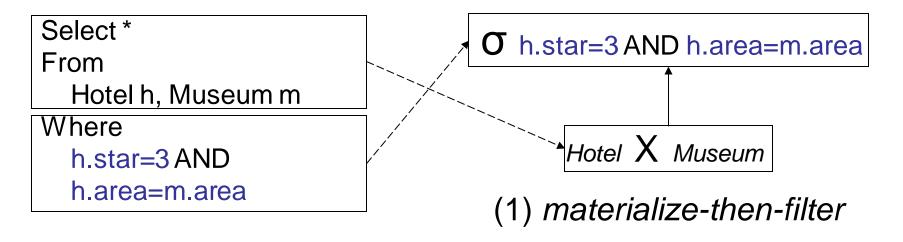


Therefore the problem is:

Unlike Boolean constructs, ranking is second class.

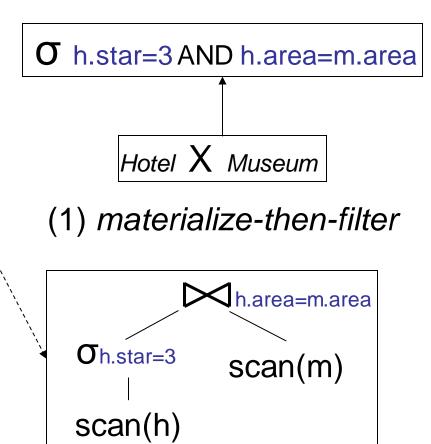
– Ranking is processed as a Monolithic component (\mathcal{R}) , always after the Boolean component (\mathcal{B}) .

How did we make Boolean "first class"?



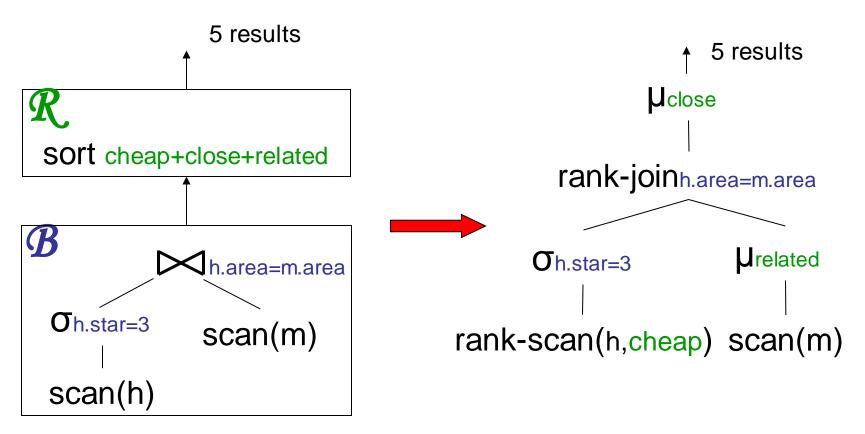
First Class: Splitting and Interleaving

Select *
From
Hotel h, Museum m
Where
h.star=3 AND
h.area=m.area



(2) **B** is *split* into *joins* and *selections*, which *interleave* with each other.

Ranking Query Plan



materialize-then-sort. naïve, overkill

split and interleave:

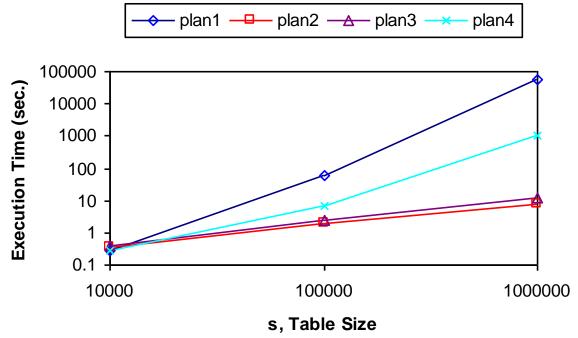
reduction of intermediate results, thus processing cost

Possibly orders of magnitude improvement

Implementation in PostgreSQL

plan1: traditional materialize-then-sort plan

plan2-4: new ranking query plans



Observations:

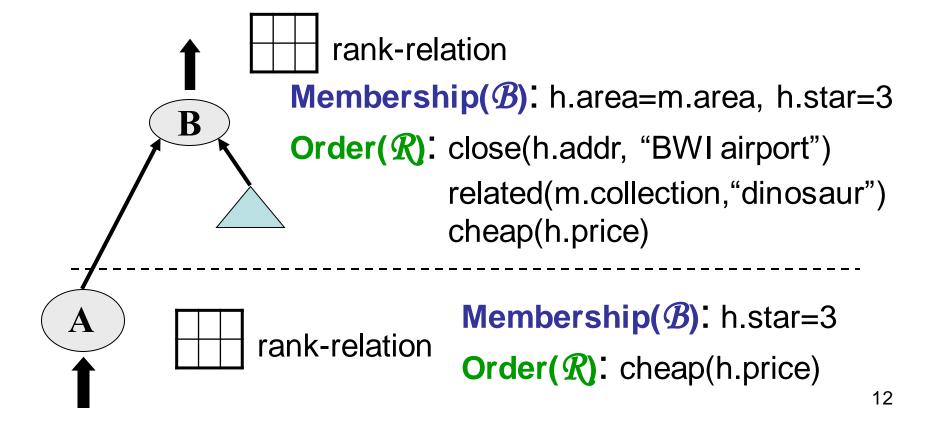
an extended plan space with plans of various costs.

RankSQL

- Goals:
 - Support ranking as a first-class query type in RDBMS;
 splitting ranking.
 - Integrate ranking with traditional Boolean query constructs.
 interleaving ranking with other operations.
- Foundation: Rank-Relational Algebra
 - data model: rank-relation
 - operators: new and augmented
 - algebraic laws
- Query engine:
 - executor: physical operator implementation
 - optimizer: plan enumeration, cost estimation

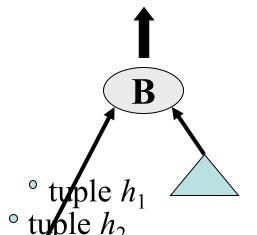
Two Logical Properties of Rank-Relation

- Membership of the tuples: evaluated Boolean predicates
- Order among the tuples: evaluated ranking predciates



Ranking Principle: what should be the order?

F=cheap + close + related



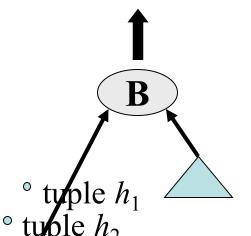
hotel	cheap	upper	museum	close	related
		bound			
h1	0.9	2.9	*	1	1
h2	0.6	2.6	*	1	1

Upper-bound determines the order:

 Without further processing h1, we cannot output any result;

Ranking Principle: upper-bound determines the order

F=cheap + close + related



hotel	cheap	upper	museum	close	related
		bound			
h1	0.9	2.9	*	1	1
h2	0.6	2.6	*	1	1
		•••	•••		

Upper-bound determines the order:

- Without further processing h1, we cannot output any result;
- Processing in the "promising" order, avoiding unnecessary processing.

Rank-Relation

• Rank-relation R_p^F

R: relation

F: monotonic scoring function over predicates (p₁, ..., p_n)

 $P \subseteq \{p_1, \ldots, p_n\}$: evaluated predicates

- Logical Properties:
 - Membership:R (as usual)
 - Order: < \forall t1, t2 \in R_P^F : t1 < t2 iff $\overline{F_P}$ [t1] < $\overline{F_P}$ [t2]. (by upper-bound)

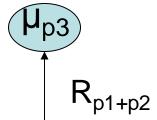
Operators

To achieve splitting and interleaving:

- New operator
 - μ : evaluate ranking predicates piece by piece. implementation: MPro (Chang et al. SIGMOD02).
- Extended operators:
 - rank-selection
 - rank-joinimplementation: HRJN (Ilyas et al. VLDB03).
 - rank-scan
 - rank-union, rank-intersection.

hotel	 p1	p2	рЗ	score
h1	0.7	0.8	0.9	2.4
h2	0.9	0.85	8.0	2.55
h3	0.5	0.45	0.75	1.7
h4	0.4	0.7	0.95	2.05
				•••

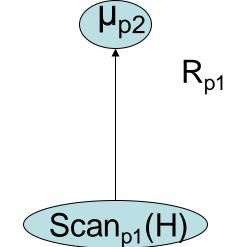
$R_{p1+p2+p3}$



hotel	upper-bound

hotel	upper-bound

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



hotel	upper-bound

hotel	 p1	<i>p</i> 2	рЗ	score
h1				
h2	0.9			
h3				
h4				

 $R_{p1+p2+p3}$

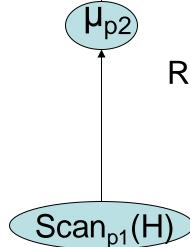


hotel upper-bound

 R_{p1+p2}

hotel	upper-bound

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



hotel upper-bound
h2

hotel	 p1	<i>p</i> 2	рЗ	score
h1	0.9	1.0	1.0	2.9
h2	0.9	1.0	1.0	2.9
h3	0.9	1.0	1.0	2.9
h4	0.9	1.0	1.0	2.9
	0.9	1.0	1.0	2.9

$R_{p1+p2+p3}$

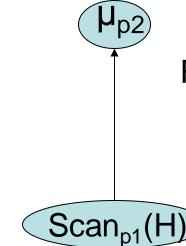


R_{p1+p2}

hotel	upper-bound

hotel	upper-bound

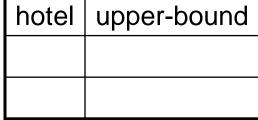
Select *
From Hotel H
Order By p1+p2+p3
Limit 1



hotel	upper-bound
h2	2.9

hotel	 p1	p2	рЗ	score
h1	0.9	1.0	1.0	2.9
h2	0.9	1.0	1.0	2.9
h3	0.9	1.0	1.0	2.9
h4	0.9	1.0	1.0	2.9
	0.9	1.0	1.0	2.9

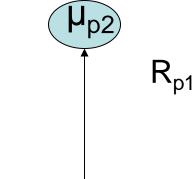
$R_{p1+p2+p3}$	r
p :	



R_{p1}	+p2
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hotel	upper-bound
h2	2.9

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



Scan_{p1}(H)

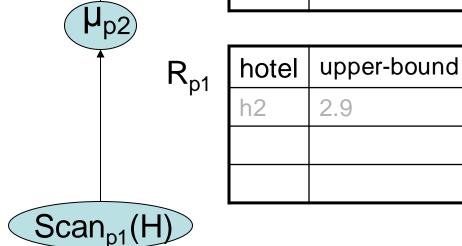
hotel	upper-bound		
h2	2.9		

hotel	 p1	p2	рЗ	score
h1	0.9	1.0	1.0	2.9
h2	0.9	0.85	1.0	2.75
h3	0.9	1.0	1.0	2.9
h4	0.9	1.0	1.0	2.9
	0.9	1.0	1.0	2.9

R _{p1+p2+p3}	hotel	upper-bound	
n3			

R.,,,,,	hotel	upper-bound	
R _{p1+p2}	h2	2.75	

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



hotel	 p1	<i>p</i> 2	рЗ	score
h1	0.7	1.0	1.0	2.7
h2	0.9	0.85	1.0	2.75
h3	0.7	1.0	1.0	2.7
h4	0.7	1.0	1.0	2.7
	0.7	1.0	1.0	2.7

R	р1+	-n2	+n:	3
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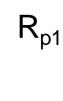
 μ_{p2}

 R_{p1+p2}

hotel	upper-bound

hotel	upper-bound
h2	2.75

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



hotel	upper-bound
h2	2.9
h1	2.7



hotel	 p1	<i>p</i> 2	рЗ	score
h1	0.7	0.8	1.0	2.5
h2	0.9	0.85	1.0	2.75
h3	0.7	1.0	1.0	2.7
h4	0.7	1.0	1.0	2.7
	0.7	1.0	1.0	2.7

$R_{p1+p2+p3}$	3
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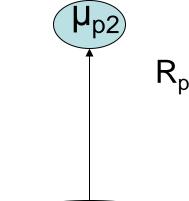


hotel	upper-bound

 R_{p1+p2}

hotel	upper-bound
h2	2.75
h1	2.5

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



Scan_{p1}(H)

hotel	upper-bound
h2	2.9
h1	2.7

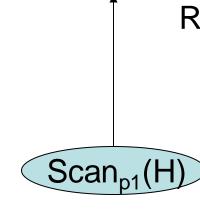
hotel	 р1	p2	рЗ	score
h1	0.7	0.8	1.0	2.5
h2	0.9	0.85	8.0	2.55
h3	0.7	1.0	1.0	2.7
h4	0.7	1.0	1.0	2.7
	0.7	1.0	1.0	2.7

$R_{p1+p2+p3}$	hotel	upper-bound
	h2	2.55
(p3)		

R_{p1+p2}	

hotel	upper-bound
h2	2.75
h1	2.5

Select *
From Hotel H
Order By p1+p2+p3
Limit 1

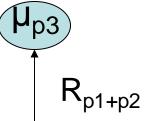


 μ_{p2}

hotel	upper-bound
h2	2.9
h1	2.7

hotel	 p1	p2	рЗ	score
h1	0.7	0.8	1.0	2.5
h2	0.9	0.85	8.0	2.55
h3	0.5	1.0	1.0	2.5
h4	0.5	1.0	1.0	2.5
	0.5	1.0	1.0	2.5

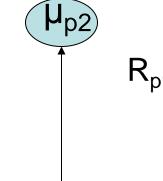
R	പ_	·p2+	Ln3
	γι +	·μΖ٦	rps



hotel	upper-bound
h2	2.55

hotel	upper-bound
h2	2.75
h1	2.5

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



Scan_{p1}(H)

hotel	upper-bound
h2	2.9
h1	2.7
h3	2.5

hotel	 p1	p2	рЗ	score
h1	0.7	0.8	1.0	2.5
h2	0.9	0.85	8.0	2.55
h3	0.5	0.45	1.0	1.95
h4	0.5	1.0	1.0	2.5
	0.5	1.0	1.0	2.5

R	n1+	-n2-	+p3
	~	P -	, 60

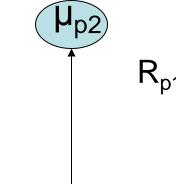


 R_{p1+p2}

hotel	upper-bound
h2	2.55

hotel	upper-bound
h2	2.75
h1	2.5
h3	1.95

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



 $Scan_{p1}(H)$

hotel	upper-bound
h2	2.9
h1	2.7
h3	2.5

hotel	 p1	<i>p</i> 2	рЗ	score
h1	0.7	0.8	0.9	2.4
h2	0.9	0.85	0.8	2.55
h3	0.5	0.45	1.0	1.95
h4	0.5	1.0	1.0	2.5
	0.5	1.0	1.0	2.5

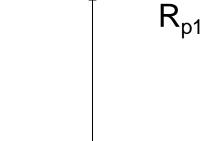


hotel	upper-bound
h2	2.55
h1	2.4

 R_{p1+p2}

hotel	upper-bound		
h2	2.75		
h1	2.5		
h3	1.95		

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



Scan_{p1}(H)

 μ_{p2}

hotel	upper-bound		
h2	2.9		
h1	2.7		
h3	2.5		

hotel	 p1	<i>p</i> 2	рЗ	score
h1	0.7	8.0	0.9	2.4
h2	0.9	0.85	8.0	2.55
h3	0.5	0.45	1.0	1.95
h4	0.5	1.0	1.0	2.5
	0.5	1.0	1.0	2.5



hotel	upper-bound	
h2	2.55	
h1	2.4	

 R_{p1+p2}

hotel	upper-bound		
h2	2.75		
h1	2.5		
h3	1.95		

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



 $Scan_{p1}(H)$

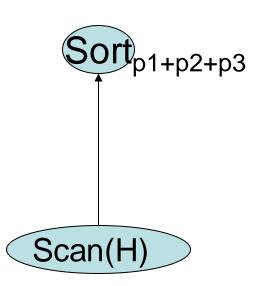
 μ_{p2}

hotel	upper-bound
h2	2.9
h1	2.7
h3	2.5

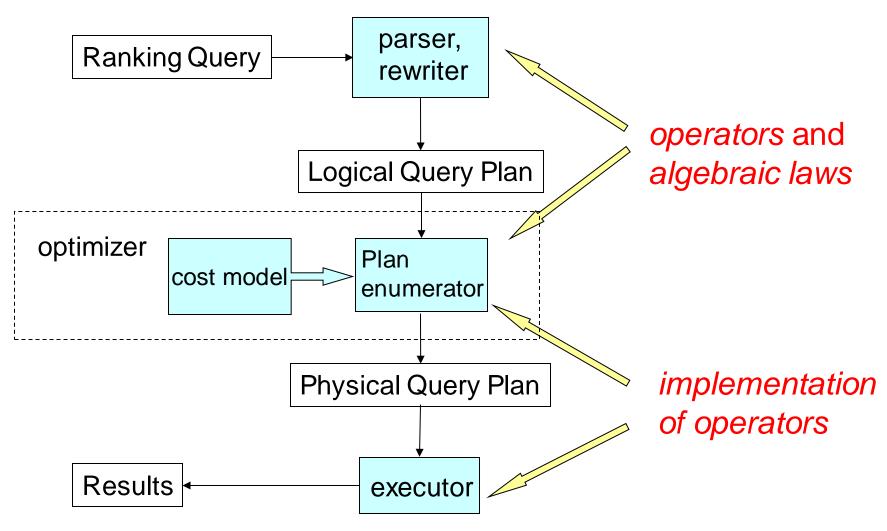
In contrast: materialize-then-sort

hotel	 р1	<i>p</i> 2	рЗ	score
h1	0.7	0.8	0.9	2.4
h2	0.9	0.85	8.0	2.55
h3	0.5	0.45	0.75	1.7
h4	0.4	0.7	0.95	2.05

Select *
From Hotel H
Order By p1+p2+p3
Limit 1



Impact of Rank-Relational Algebra



Optimization

Two-dimensional enumeration:
 ranking (ranking predicate scheduling)
 and
 filtering (join order selection)

Sampling-based cardinality estimation

Two-Dimensional Enumeration

- (1 table, 0 predicate)
 seqScan(H), idxScan(H), seqScan(M), ...
- (1 table, 1 predicate) $\text{rankScan}_{\text{cheap}}(H), \quad \mu_{\text{cheap}}(\text{seqScan}(H)), \dots$
- (1 table, 2 predicates) $\mu_{close}(rankScan_{cheap}(H)), ...$
- (2 table, 0 predicate)
 NestLoop(seqScan(H), seqScan(M)), ...
- (2 table, 1 predicate)
 NRJN(rankScan_{cheap}(H), seqScan(M)),...
- and so on...

Related Work

Middleware

Fagin et al. (PODS 96,01), Nepal et al. (ICDE 99), Günter et al. (VLDB 00), Bruno et al. (ICDE 02), Chang et al. (SIGMOD 02)

RDBMS, outside the core

Chaudhuri et al. (VLDB 99), Chang et al. (SIGMOD 00), Hristidis et al. (SIGMOD 01), Tsaparas et al. (ICDE 03), Yi et al. (ICDE 03)

- RDBMS, in the query engine
 - Physical operators and physical properties
 Carey et al. (SIGMOD 97), Ilyas et al. (VLDB 02, 03, SIGMOD 04), Natsev et al. (VLDB 01)
 - Algebra framework
 Chaudhuri et al. (CIDR 05)

Conclusion: RankSQL System

Goal:

Support ranking as a first-class query type; Integrate ranking with Boolean query constructs.

Our approach:

Algebra: rank-relation,

new and augmented rank-aware operators,

algebraic laws

- Optimizer: two-dimensional enumeration,

sampling-based cost estimation

Implementation: in PostgreSQL

Welcome to our demo in VLDB05!