CSE 5311 Test 1 - Closed Book Spring 2019	Name	ame as on your UTA ID card
Multiple Choice/Short Answer:		
<ol> <li>Write your answer, letter or value, on</li> <li>CIRCLED ANSWERS DO NOT CO</li> <li>2 points each</li> </ol>	n the line () to the LE OUNT.	FT of each problem.
1. The number of potential probe sequences prime) is:	s when using double hashing	with a table with $m$ entries ( $m$ is
A. $O(\log m)$ B. $m$	C. $m(m-1)$ D. $m!$	
A. INIT had an amortized cost of 1.  B. SUCC had an actual cost determine C. SUCC had an amortized cost of 2.  D. The potential was defined with reg	ed by the number of edges fo	llowed.
Suppose you already have 10 different counumber of boxes for obtaining a coupon described in the coupo	upons when there are 20 cou	non types. What is the arms of 1
← Write your answer on the line	20	
. Mapping a multi-way tree to a binary tree priority queue approach?	(left child, right sibling repre	esentation) is part of which
A. binary heaps B. binomial hear	ps C. leftist heaps	D. pairing heaps
The balance factors in an AVL tree are con	nputed as:	
A. height <sub>right</sub> - height <sub>left</sub> B. same as the null path length in a left C. the difference of the number of node D. the distance from a node to the root	es in the left and right subtre	ees
How many inversions are there for the lists	3, 1, 2, 4, 5 and 2, 5, 4, 3, 1	9 3 > 3

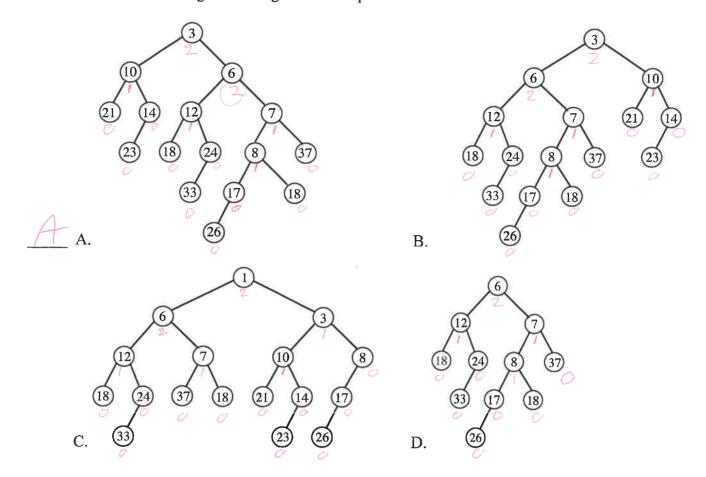
7. Assuming a random *n*-permutation is provided, the expected number of hires for the hiring problem is:

 $\triangle$  A.  $H_n$  B. 2 C.  $\sqrt{n}$  D.  $\ln \ln n$ 

← Write your answer on the line

5

## 8. Which of the following is not a legal leftist heap?



## 9. Which property does not hold for binomial heaps?



- A. Performing n INSERT operations into an empty heap will take O(n) time.
- B. The number of trees is based on the binary representation of the number of stored items.
- C. Decrease-Key takes O(1) time.  $\times$
- D. MINIMUM takes  $O(\log n)$  time.

## 10. The perfect hashing method discussed in class depends on which fact?



A. 
$$\sum_{k=0}^{\infty} x^k = \frac{1}{1-x}$$
  $0 < x < 1$ 

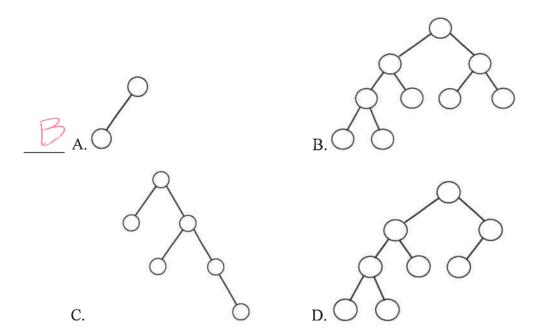
- B.  $\ln n < H_n < \ln n + 1$
- C. The expected number of probes for a successful search in Brent's method is less than 2.5.
- D. The probability of collisions among n keys stored in a hash table of size  $n^2$  is less than 0.5.
- 11. If a Fibonacci tree appears as a subtree of an AVL tree, which nodes would be assigned a balance factor of 0?



- A. none of them
- B. only the leaves
- C. only the root
- D. the leaves and the root

12. What is the worst-case number of	rotations when	n perfo	ming del	etion on a	ın AV	VL tre	e?	-
$\triangle$ A. $\Theta(n)$ B. $\Theta(1)$	C. $\Theta(\log n)$	D. No	rotations	are ever	used			
13. The master method summarizes th					-	ie.		
A. asymptotic analysis	B. potential fur	nctions	C. recur	sion trees	Ι	D. su	bstitution metho	od
14. Which of the following is not true	regarding Blo	om filte	rs?					
A. They are an especially con B. Several hash functions are C. The optimal bit array size D. An indication that a candid	used.  depends on the	e numbe	r of item	s and the f	false rrect.	posit	ive probability.	
15. When are Fibonacci trees used?								
A. Constructing a priority que B. Demonstrating worst-case C. Defining the potential func D. Demonstrating worst-case	behaviors for re tion for splay t	red-blac trees.	k trees.	mplexity 1	for D	ECRE	ASE-KEY.	
16. When performing selection in wor medians are computed in the first r	st-cast linear ti ound?	ime for	n numbei	s, roughly	y how	v mar	ıy column	
A. $\frac{n}{5}$ B. $m$ , the	e median-of-m	nedians	C	0.7 <i>n</i>	D	). W	$\left(\frac{n}{5}\right)$	
17. The main difference between MTF	and OPT for s	self-org	anizing li	near lists	is:			
A. OPT can do transpositions								
B. OPT is given the entire requartime	iest sequence i	in advaı	ce, while	MTF rec	eives	s the	requests one-at-	
<ul><li>C. MTF is given the entire req</li><li>a-time </li><li>D. MTF counts inversions </li></ul>	uest sequence i	in adva	nce, while	e OPT rec	eives	s the	requests one-at-	
18. When of the following is not an app	olication of uni	ion-find	trees?					
A. Finding the median of a set B. Kruskal's algorithm								
C. Off-line least common ancer D. Maximum cardinality <i>k</i> -cold		of interv	als					
19. What is required when calling union	on(i,j) for	maintai	ning disj	oint subse	ts?			
A. i and j are roots for difference.  C. i and j are roots for the same				ancestor are in the			e of the trees	

20. Which of the following binary trees has multiple legal colorings as a red-black tree?



21. When performing selection in worst-cast linear time, roughly what fraction of the set of n keys is kept (in the worst case) for the next round?



- A. 20%
- B. 30%
- C. 70%
- D. 80%

22. When using Brent's rehash, the number of previously inserted keys that may move is:



- A. 1 B. 2
- C.  $\frac{1}{a}$  D.  $H_m$ , where m is the number of stored keys

23. Which priority queue implementation generalizes binary heaps by increasing the branching?



- A. Binomial heaps B. d-heaps C. Pairing heaps
- D. Leftist heaps

24. What is minimized in the dynamic programming solution to the subset sum problem?



- A. The number of input values used to sum to each C(i)

C. The index stored for each C(i)

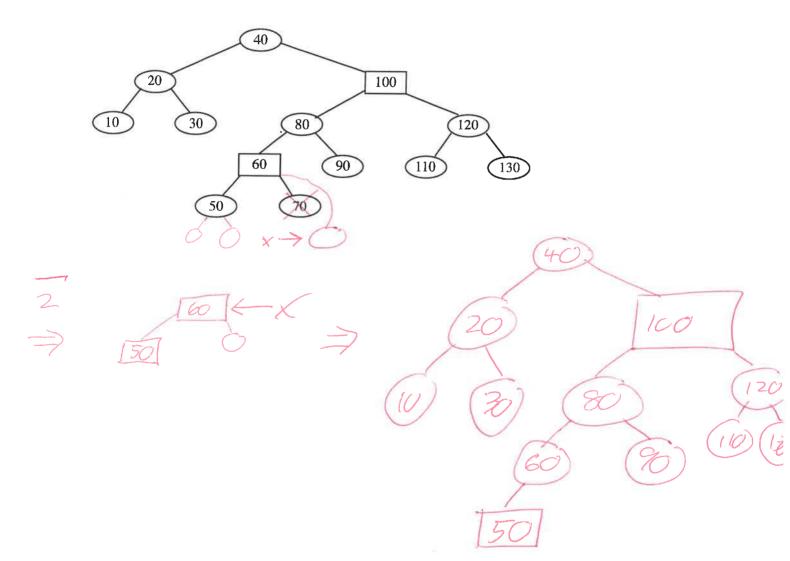
D. m

25. The recursion tree for mergesort has which property?



- A. it leads to a definite geometric sum B. it leads to a harmonic sum
- C. each level has the same contribution D. it leads to an indefinite geometric sum

1. Show the result of deleting 70 from this red-black tree (preserving red-black properties). Be sure to indicate the cases that were used. (10 points)



2. The hash table below was created using double hashing with Brent's rehash. The initial slot  $(h_1(key))$  and rehashing increment  $(h_2(key))$  are given for each key. Show the result from inserting 284 using Brent's rehash when  $h_1(284) = 4$  and  $h_2(284) = 3$ . (10 points)

	key	$h_1(key)$	$h_2(key)$	
0	171	3	4	041
1				
2	912	2	1	23456
3	801	3	4	3041
94	102	4	1	4.56
5	212	2	3	5 1
6	W. L			
	key		-	284 4036 Neuprobes 1234 Oldenbes 2230 £ (3)464
0	171			
1		3		
2	912			
3	801			
4	284			
5	212	-		
6	102			

Indicate the additional number of probes for searching for all keys once after 284 is inserted:



3. Construct the final optimal binary search tree (using Knuth's root trick) and give its cost. SHOW YOUR WORK. (10 points)

```
n=6
                      w[2][3]=0.150000
                                                              c(1,1) cost 0.000000
q[0]=0.03;
                      w[2][4]=0.280000
                                                              c(2,2) cost 0.000000
key[1]=1;
                      w[2][5]=0.480000
                                                              c(3,3) cost 0.000000
p[1]=0.20;
                      w[2][6]=0.610000
                                                              c(4,4) cost 0.000000
q[1]=0.01;
                      w[3][3]=0.020000
                                                              c(5,5) cost 0.000000
key[2]=2;
                      w[3][4]=0.150000
                                                              c(6,6) cost 0.000000
p[2]=0.15;
                      w[3][5]=0.350000
                                                              c(0,1) cost 0.240000 1
                                                              c(1,2) cost 0.190000 2
q[2]=0.03;
                      w[3][6]=0.480000
                      w[4][4]=0.030000
                                                              c(2,3) cost 0.150000 3
key[3]=3;
p[3]=0.10;
                     w[4][5]=0.230000
                                                              c(3,4) cost 0.150000 4
q[3]=0.02;
                     w[4][6]=0.360000
                                                              c(4,5) cost 0.230000 5
                      w[5][5]=0.050000
                                                              c(5,6) cost 0.180000 6
key[4]=4;
                     w[5][6]=0.180000
                                                              c(0,2) cost 0.610000 1(,2)
p[4]=0.10;
                                                              c(1,3) cost 0.460000 2(,3)
q[4]=0.03;
                     w[6][6]=0.030000
key[5]=5;
                      Building c(0,2) using roots 1 thru 2
                                                              c(2,4) cost 0.430000 4(3,)
                     Building c(1,3) using roots 2 thru 3
                                                              c(3,5) cost 0.500000 5(4,)
p[5]=0.15;
q[5]=0.05;
                     Building c(2,4) using roots 3 thru 4
                                                              c(4,6) cost 0.540000 5(,6)
key[6]=6;
                     Building c(3,5) using roots 4 thru 5
                                                              c(0,3) cost 0.930000 2(1,3)
p[6]=0.10;
                     Building c(4,6) using roots 5 thru 6
                                                              c(1,4) cost 0.780000 3(2,4)
                     Building c(0,3) using roots 1 thru 2
q[6]=0.03;
                                                              c(2,5) cost 0.860000 4(3,5)
                     Building c(1,4) using roots 2 thru 4
w[0][0]=0.030000
                                                             c(3,6) cost 0.810000 5(4,6)
w[0][1]=0.240000
                     Building c(2,5) using roots 4 thru 5
                                                             c(0,4) cost 1.340000 2(1,4(3,))
w[0][2]=0.420000
                     Building c(3,6) using roots 5 thru 5
                                                             c(1,5) cost 1.330000 3(2,5(4,))
w[0][3]=0.540000
                     Building c(0,4) using roots 2 thru 3
                                                             c(2,6) cost 1.220000 5(4(3,),6)
                     Building c(1,5) using roots 3 thru 4
                                                             c(0,5) cost 1.970000 (2(1,4(3,5))
w[0][4]=0.670000
                     Building c(2,6) using roots 4 thru 5
                                                             c(1,6) cost 1.730000 5(3(2,4),6)
w[0][5]=0.870000
                                                             c(0,6) cost ??? ????????????????
                     Building c(0,5) using roots 2 thru 3
w[0][6]=1.000000
w[1][1]=0.010000
                     Building c(1,6) using roots 3 thru 5
                                                              3(1(,2),5(4,6)
w[1][2]=0.190000
                     Building c(0,6) using roots ? thru ?
                     Counts - root trick 33 without root
w[1][3]=0.310000
w[1][4]=0.440000
                         trick 50
w[1][5]=0.640000
                     Average probe length is ???
w[1][6]=0.770000
                     trees in parenthesized prefix
w[2][2]=0.030000
                     c(0,0) cost 0.000000
        1.0 + c(0,1) + c(2,6)
        1.0 + c(0,2) + c(3,6) =
        1.0 + C(0,3) + C(4,6) =
       1.0 + C(0,4) + C(5,6) = 2.52
```

4. Use the substitution method to show that  $T(n) = 3T(\frac{n}{3}) + n$  is in  $O(n \log n)$ . (You do not need to show that T(n) is in  $O(n \log n)$ .) (10 points)

Suppose  $T(K) \le CK \log_3 K \le K < n$   $T(\frac{n}{3}) \le C \frac{n}{3} \log_3 \frac{n}{3}$   $= C \frac{n}{3} \log_3 n - C \frac{n}{3}$   $T(n) \le 3 \left[ C \frac{n}{3} \log_3 n - C \frac{n}{3} \right] + n$   $= cn \log_3 n - cn + n$   $\le cn \log_3 n - cn + n$   $\le cn \log_3 n - cn + n$ 

5. Give the range of possible heights for an AVL tree with 100 keys. Your answer should be two natural numbers giving the minimum and maximum heights. (A tree with one node has height 0.) Show your work! (10 points)

Complete tree of height
$$h \quad has \quad 2^{h+1} \quad nodes$$

$$h=5 \Rightarrow 2^{6}-1=63$$

$$h=6 \Rightarrow 2^{7}-1=127$$

maximum 8

F: b	Tree5			
height	nade5			
0	2			
7	4			
2 3	7			
4	12			
5	20			
6	33			
7	54			
> 8	88			
9	(143)	5.nce smaller	100 is	be
				7