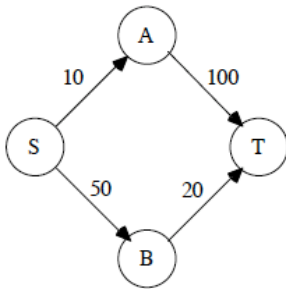


## CSE 5311: Homework 2

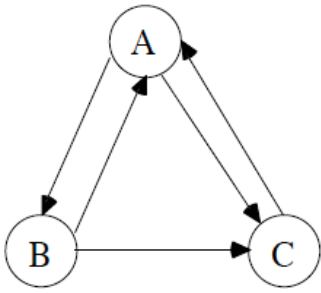
1. 26.2-3, but also include an edge of capacity 10 from  $v_1$  to  $v_2$ . Also solve using preflow-push
2. Determine the lattice of solutions for the following stable marriage problem, given the following precedence lists ordered from most favored to least favored:

Men				Women			
1	2	3	4	1	2	3	4
1	2	3	4	4	3	2	1
2	1	4	3	3	4	1	2
3	4	1	2	2	1	4	3
4	3	2	1	1	2	3	4

3. 22-1
4. Find a maximum flow in the following network using the preflow-push algorithm.



5. Give the Knuth-Morris-Pratt fail links (both methods) and the Z algorithm prefix lengths for the search pattern `abracadabra`.
6. Suppose that matrix multiplication is implemented in a recursive, decomposition fashion like Strassen's method. However, instead of using his equations we use the everyday ones, i.e.  $c_{ij} = a_{i1} * b_{1j} + a_{i2} * b_{2j}$ . What is the asymptotic complexity, based on the number of scalar multiplies and additions/subtractions?
7. It is easy to test in polynomial time if the sum of the lengths for a set of line segments  $\leq$  a given value  $k$ . This problem may be generalized to two dimensions to give an NP-complete problem. In the generalized problem, the set of line segments is replaced by a set of rectangles (with dimensions in  $x$  and  $y$ ). We would like to test if the rectangles will fit (without overlap, but they may touch) within another given rectangle, but the given rectangles must remain oriented to the  $x$  and  $y$  axes, i.e. do not rotate them. Prove that this problem is NP-complete by a simple reduction.
8. In many cases, transitivities may be removed from a directed graph. Instead of allowing you to remove the transitivities you please until none remain, you are required to strategically remove as many as possible for that graph, i.e. removing a transitive edge may make other previously transitive edges non-transitive. As an example, consider:



If I remove edge (B,C), no more transivities remain, but I could also remove both (B,A) and (A,C) to leave no transivities.

Prove: Testing if there are k edges that may be removed from a directed graph while still preserving all directed paths from the original graph is NP-complete.

Hint: Use an extraordinarily simple reduction from Directed Hamiltonian Circuit. You may assume that the DHC graph has a directed path between each pair of vertices.

9. 34-1

10. 35.2-2

11. 35.2-3

12. 35.2-4

13. 35.3-1

14. 35.3-2