

0-1 Knapsack Problem

Given _____, _____, and arrays _____, _____

Find _____ (rarely _____)

Have _____ of each item.

(_____, _____), (_____, _____), (_____, _____), (_____, _____)

How do we solve it?

1. _____

2. _____

3. _____

Dynamic Programming (DP) Solution

Solve all smaller problems (from problem size 0 to current problem size) => table

Observations:

- Does order of picked items matter? (e.g. 1,2,3 vs 2,1,3?) _____
- _____
- What is a problem of size 0? _____
- What is a problem of size 1? Can I have a single problem of size 1 (regarding items)?

Implementation:

row/column indexes = 0

item's weight > current weight (column idx) = value of cell 1 row above current cell

item's weight ≤ current weight (column idx) = max between value of cell 1 row above current cell & value of current item + value of cell 1 row above it and current weight - current item's weight columns before current cell

Equations:

sol = matrix/2D array storing answers for (all) smaller pbs TC:

$w[i]$ = weight of current item

SC:

$v[i]$ = value of current item

k = current weight (column indexes)

$sol[0][k] = \underline{\hspace{2cm}}$, for $\underline{\hspace{2cm}}$

$sol[i][0] = \underline{\hspace{2cm}}$, for $\underline{\hspace{2cm}}$

$sol[i][k]$:

If $k < w[i]$: $\underline{\hspace{10cm}}$

If $k \geq w[i]$: $\underline{\hspace{10cm}}$

ID	Weight	Value
1		
2		
3		
4		

Max Weight =

sol table below. Indicate in table picked item with "*" and not picked with "."

$sol[0][5] = \underline{\hspace{10cm}}$

$sol[1][4] = \underline{\hspace{10cm}}$

$sol[3][5] = \underline{\hspace{10cm}}$

$sol[4][8] = \underline{\hspace{10cm}}$

$sol[3][5] = \max \{ \underline{\hspace{5cm}}, \underline{\hspace{5cm}} \}$

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16

Backtracking:

Final Value Achieved: _____ Items Picked: _____

DP solution variations:

- Space saving:

Knapsack problem variations:

1. _____
2. _____
3. _____

```
/*Arrays v and w have info from index 1: first item has value v[1] and weight w[1]
*/
```

```
int knapsack01(int W, int n, int * v, int * w){
    int sol[n+1][W+1];
    for(k=0; k<=W; k++) { sol[0][k] = 0;}
    for(i=1; i<=n; i++) {
        for(k=0;k<=W;k++) {
            sol[i][k] = sol[i-1][k]; // solution without item i
            if (k>w[i]) {
                with_i = v[i]+sol[i-1][k-w[i]];
                if (sol[i][k] < with_i) { // better choice
                    sol[i][k] = with_i; // keep it
                }
            }
        }
    }
    return sol[n][W];
}
```

} // Time: $\Theta(\rule{1cm}{0.4pt})$ Space: $\Theta(\rule{1cm}{0.4pt})$ pseudo polynomial in W

// need $\Theta(n)$ bits to store n items (values and weights) , but only $\log_2(W)$ bits to store value W

Greedy:

Max Weight =		
ID	Weight	Value
1		
2		
3		
4		

Criterion: _____

Work:

[illegible]

Criterion: _____

Work:

[illegible]

Optimality:

0/1 Knapsack:

Fractional Knapsack:

Other variations:

Job Scheduling:

Criteria	Max Total Value/Profit	Max Number of Jobs
Max Value/Length		
Max Value		
Min Length		
Finishes Last		
Starts First		
Finishes First		
Starts Last		

Difference between Greedy and Dynamic Programming:

Greedy:

Dynamic Programming

Brute Force

Idea:

How many possible combinations of items are there for N items? _____

How to generate all possible combinations of items?

- Idea
- TC