

# Asymptotic Bounds

## Review:

Use ratio and limit to infinity to compare function growth. Examples:

$$a) \quad N^2 < N^3 \quad \frac{N^2}{N^3} = \frac{1}{N} \quad \text{use } 1 \text{ grows slower than } N \quad \text{or} \quad \lim_{N \rightarrow \infty} \frac{1}{N} = 0 \quad \Rightarrow \text{top grows slower}$$

$$b) \quad N\sqrt{N} > N \quad \frac{N\sqrt{N}}{N} = \frac{\sqrt{N}}{1} \quad \text{use } \sqrt{N} \text{ grows faster than } 1 \quad \text{or} \quad \lim_{N \rightarrow \infty} \frac{\sqrt{N}}{1} = \infty \quad \Rightarrow \text{top grows faster}$$

$$c) \quad N\sqrt{N} < N^2 \quad \frac{N\sqrt{N}}{N^2} = \frac{1}{\sqrt{N}} \quad \text{use } 1 \text{ grows slower than } \sqrt{N} \quad \text{or} \quad \lim_{N \rightarrow \infty} \frac{1}{\sqrt{N}} = 0 \quad \Rightarrow \text{top grows slower}$$

$$d) \quad \lg N \approx \log_3 N \quad \frac{\lg N}{\log_3 N} = \frac{\lg N}{\frac{\lg N}{\lg 3}} = \frac{\lg N}{1} * \frac{\lg 3}{\lg N} = \lg 3 \quad \text{use } \lg 3 \text{ is a constant (no } N) \quad \text{or} \quad \lim_{N \rightarrow \infty} \frac{\lg N}{\log_3 N} = \lg 3 \quad \Rightarrow \text{same growth}$$

Examples of functions that have  $N^2$  growth:

Speed of Growth:

\_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_ < \_\_\_\_\_

Fill in each column functions that have the same growth as the column label.

1	$\lg N$	$N$	$N \lg N$	$N^2$	$N^3$	$2^N$

The functions listed in a column are \_\_\_\_\_ of the function that labels the column. E.g. \_\_\_\_\_ = \_\_\_\_\_ ( $N^2$ )

From ordering by growth to asymptotic notation

.....  
 .....

Symbol:  $\Theta$   $O$   $o$   $\Omega$   $\omega$

Meaning:

Examples:  $N^3/10 - 500N^2 - 1000 = O(N^3)$  True/False Solution: find dominant term(s) and compare their growth

Notation abuse: = instead of  $\in$

Finding  $\Theta$ ,  $O$ ,  $\Omega$  for an algorithm.

Worst case is $\Theta(\underline{\hspace{1cm}})$ then the algorithm is $\underline{\hspace{1cm}}$ ( $\underline{\hspace{1cm}}$ )	Insertion sort: Worst case is $\Theta(\underline{\hspace{1cm}})$ then insertion sort is $\underline{\hspace{1cm}}$ ( $\underline{\hspace{1cm}}$ )
Best case is $\Theta(\underline{\hspace{1cm}})$ then the algorithm is $\underline{\hspace{1cm}}$ ( $\underline{\hspace{1cm}}$ )	Best case is $\Theta(\underline{\hspace{1cm}})$ then insertion sort is $\underline{\hspace{1cm}}$ ( $\underline{\hspace{1cm}}$ )

An algorithm is  $\Theta(\underline{\hspace{1cm}})$  iff best case and worst case have  $\underline{\hspace{1cm}}$  TC.

Exercise: Assume:

- Alg 1 is  $O(N^2)$  .....can be: ..... cannot be: .....
- Alg 2 is  $\Theta(N \lg N)$  can be: ..... cannot be: .....
- Alg 3 is  $\Omega(N)$  ..... can be: ..... cannot be: .....
- what could be the TC function in each case?
- which one is “better”? (define better)

Symbol	Name	Meaning	Notation	Examples	Limit theorem
$\omega$					
$\Omega$					
$\Theta$					
$O$					
$o$					

Ex: Fill in s.t the statements are correct

$TC = \Theta(\log_2(N)) \Rightarrow TC = o(\underline{\hspace{1cm}}), TC = O(\underline{\hspace{1cm}}), TC = \Omega(\underline{\hspace{1cm}}), TC = \omega(\underline{\hspace{1cm}})$

$\omega: \underline{\hspace{1cm}} \quad O: \underline{\hspace{1cm}} \quad \Omega: \underline{\hspace{1cm}} \quad o: \underline{\hspace{1cm}}$

$TC = \Theta(N^2) \Rightarrow TC = o(\underline{\hspace{1cm}}), TC = O(\underline{\hspace{1cm}}), TC = \Omega(\underline{\hspace{1cm}}), TC = \omega(\underline{\hspace{1cm}})$

$\omega: \underline{\hspace{1cm}} \quad O: \underline{\hspace{1cm}} \quad \Omega: \underline{\hspace{1cm}} \quad o: \underline{\hspace{1cm}}$

$TC = O(N^2)$

$\omega: \underline{\hspace{1cm}} \quad O: \underline{\hspace{1cm}} \quad \Omega: \underline{\hspace{1cm}} \quad o: \underline{\hspace{1cm}}$

### Asymptotic Bounds as Limits:

$$f(n) = \omega(g(n)) \Leftrightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$$

English Translation:

$$f(n) = \Omega(g(n)) \Leftrightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \text{ or } c \neq 0$$

English Translation:

$$f(n) = \Theta(g(n)) \Leftrightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = c \neq 0 \quad (\text{limit is a non-zero constant})$$

English Translation:

$$f(n) = O(g(n)) \Leftrightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \text{ or } c$$

English Translation:

$$f(n) = o(g(n)) \Leftrightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$$

English Translation:

### Properties:

1.

2.

3.

4.

5.

\*\*\* Transitivity (From Discrete Structures): If  $a > b$  and  $b > c$ , then  $a > c$ . This concept can be applied to TC bounds as well.

6.

7.