

Finding Time Complexity for Recurrences:

Identify the number of times a recursive call _____ and what the new _____ is
Local time complexity is _____

Note:

- c is generally used as a _____
- c _____ $\Theta(1)$
- n _____ cn _____ $\Theta(n)$

```
int foo(int N){
    int a,b,c;
    if(N<=3) return 1500; // Note N<=3
    a = 2*foo(N-1);
    // a = foo(N-1)+foo(N-1);
    printf("A");
    b = foo(N/2);
    c = foo(N-1);
    return a+b+c;
}
```

Base case: $T(\underline{\hspace{1cm}}) = \underline{\hspace{1cm}}$

Recursive case: $T(\underline{\hspace{1cm}}) = \underline{\hspace{1cm}}$

$T(N)$ gives us the Time Complexity for $\text{foo}(N)$. We need to solve it (find the closed form)

```
void bar(int N){
    int i,k,t;
    if(N<=1) return;
    bar(N/5);
    for(i=1;i<=5;i++){
        bar(N/5);
    }
    for(i=1;i<=N;i++){
        for(k=N;k>=1;k--){
            for(t=2;t<2*N;t=t+2)
                printf("B");
        }
    }
    bar(N/5);
}
```

Base case: $T(\underline{\hspace{1cm}}) = \underline{\hspace{1cm}}$

Recursive case: $T(\underline{\hspace{1cm}}) = \underline{\hspace{1cm}}$

Solve $T(N)$

Let N be the number of elements to process in this call. $N = \text{right} - \text{left} + 1$

```
int binary_search(int A[], int left, int right, int v){
    int m = left+(right-left)/2;
    if (left > right) return -1;
    if (v == A[m]) return m;
    if (v < A[m])
        return binary_search(A, left, m-1, v);
    else
        return binary_search(A, m+1, right, v);
}
```

Recurrence: base case: $T(\quad) = T(\quad) = \underline{\hspace{2cm}}$

recursive case: $T(\quad) = \underline{\hspace{2cm}}$

Draw TC tree. Use it to find TC. TC =

SC =

```
Merge_sort(A, le, r) //N = ri-le+1
    if (le>=ri) return
    else
        m = floor(le+(ri-le)/2)
        Merge_sort(A, le, m);
        Merge_sort(A, m+1, ri);
        Merge(A, le, m, ri);
```

Recurrence: base case: $T(\quad) = T(\quad) = \underline{\hspace{2cm}}$

recursive case: $T(\quad) = \underline{\hspace{2cm}}$

```
Merge(A, le, m, ri)
1  n1=m-le+1+1 // +1 for inf
2  n2=ri-m+1 // +1 for inf
3  let L[n1], R[n2] be arrays
4  for j=0 to n1-2
5      L[j]=A[le+j]
6  for j=0 to n2-2
7      R[j]=A[m+1+j]
8  L[n1] = inf
9  R[n2] = inf
10 j=0,
11 i=0
12 for k=le to ri
13     if L[i] ≤ R[i]
14         A[k]=L[i]
15         i++
16     else
17         A[k] = R[j]
18         j++
// pseudocode
// - indentation => instruction group in {}
// - loops: for k=le to ri means for(k=le; k<=ri; k++)
```