

# Trees

## (Part 1, Theoretical)

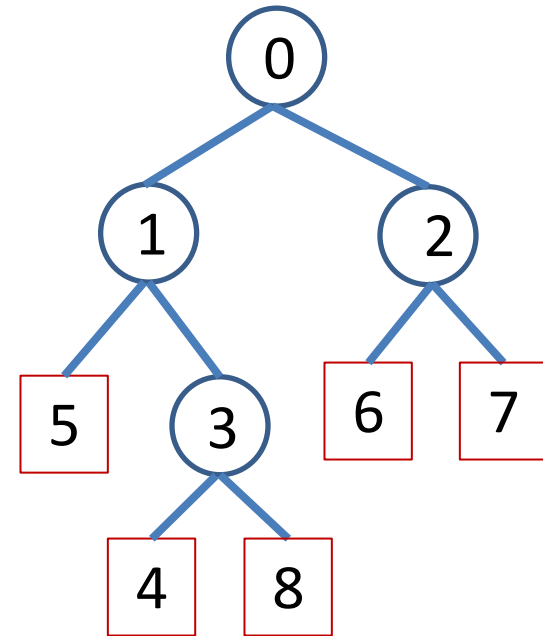
CSE 3318 – Algorithms and Data Structures  
University of Texas at Arlington

# Trees

- Trees are a natural data structure for representing specific data.
  - Family trees.
  - Organizational chart of a corporation, showing who supervises who.
  - Folder (directory) structure on a hard drive.
- Theoretical usage – analysis of recursive functions with 2 or more recursive calls

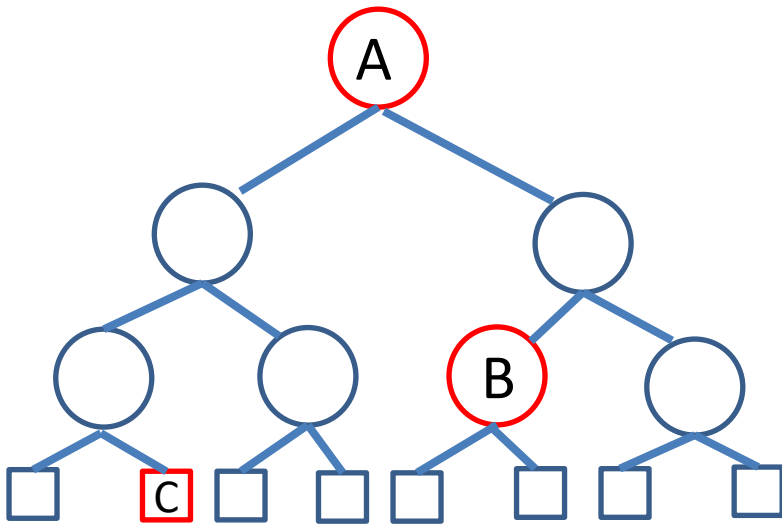
# Terminology

- **Root:** 0 (has no parent)
- **Path:** 0-2-6, 1-3, 1-3-4
- **Parent vs child**
- Ascendants vs descendants:
  - ascendants of 3: 1, 0
  - descendants of 1: 5, 3, 4, 8
- **Internal nodes:** 0, 1, 2, 3
  - Have 1 or more children
- **Leaves:** 5,4,6,7,8
  - Have no children
- **Subtree**



# Terminology - Worksheet

- The **level** of the root is defined to be 0.
- The **level** of each node is defined to be 1+ the level of its parent.
- The **depth** of a node is the number of edges from the root to the node.  
(It is equal to the level of that node.)
- The **height** of a node is the number of edges *from the node to the deepest leaf*.  
(Treat that node as the root of a small tree)



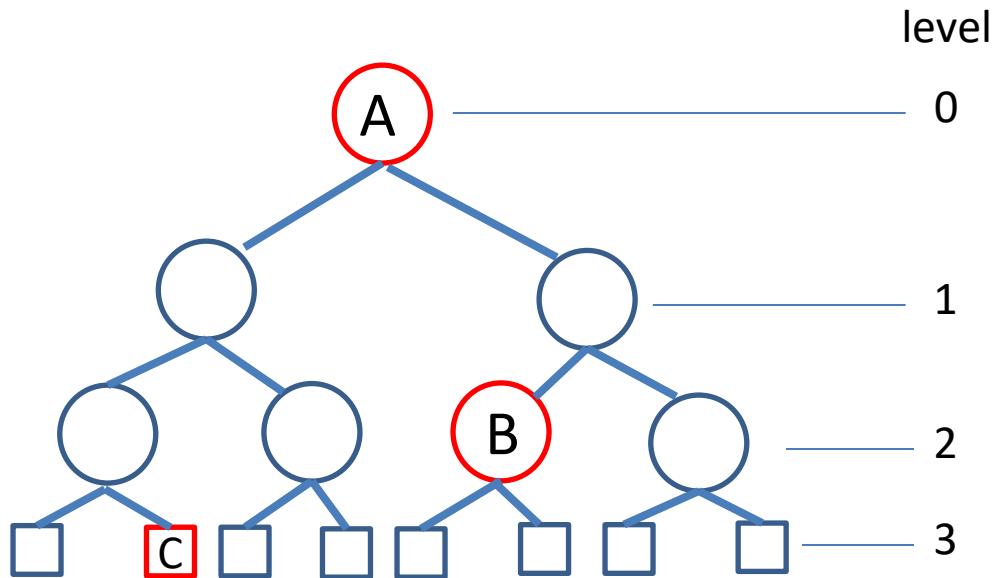
Node	level	depth	height
A			
B			
C			

Practice:

- Give the level, depth and height for each of the red nodes.
- How many nodes are on each level? \_\_\_\_\_

# Terminology - Answers

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Node	level	depth	height
A	0	0	3
B	2	2	1
C	3	3	0

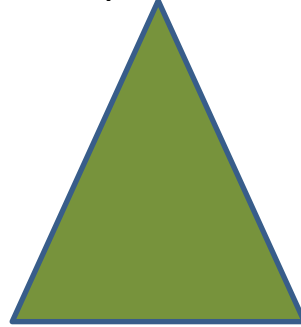
Practice:

- Give the level, depth and height for each of the red nodes.
- How many nodes are on each level? 1, 2, 4, 8

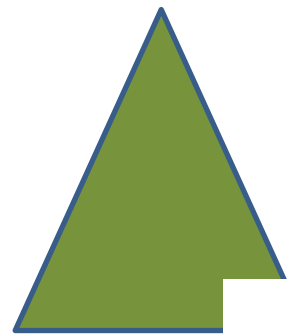
# Types of binary trees

- **Complete (or perfect)**– each internal node has exactly 2 children and all the leaves are on the same level.
  - E.g. ancestry tree (anyone will have exactly 2 parents).
- **Full [binary]** – every node has exactly 0 or 2 children.
  - Tree of recursive calls when there are 2 recursive calls
    - E.g. tree generated by the Fibonacci recursive calls.
    - => *his properties are used in analyzing time complexity of such rec fcts.*
  - *Binary tree.*
- **Nearly complete tree** – every level, except for possibly the last one is completely filled and on the last level, all the nodes are as far on the left as possible.
  - E.g. the heap tree.
  - Height:  $\lfloor \lg N \rfloor$  and it can be stored as an array.

Complete tree



Nearly Complete tree



# Complete Binary Trees

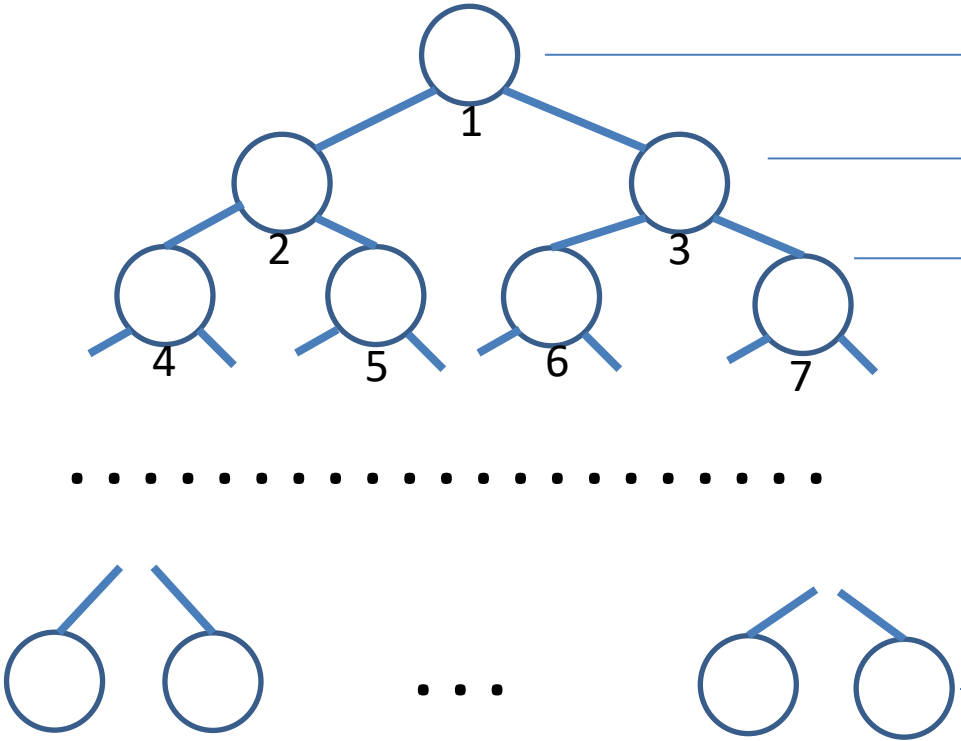
A **complete binary tree** with  $N$  nodes and height  $h$ , has:

- $\lceil N/2 \rceil$  leaves (half the nodes are on the last level)
- $\lfloor N/2 \rfloor$  internal nodes (half the nodes are internal)
- Height :  $h = \lfloor \lg N \rfloor$
- Levels :  $\lfloor \lg N \rfloor + 1$  ( $= \lg(N+1)$  )

$$\sum_{k=0}^h 2^k = 2^{h+1} - 1$$

In the other direction:  $N = 2^{h+1} - 1$

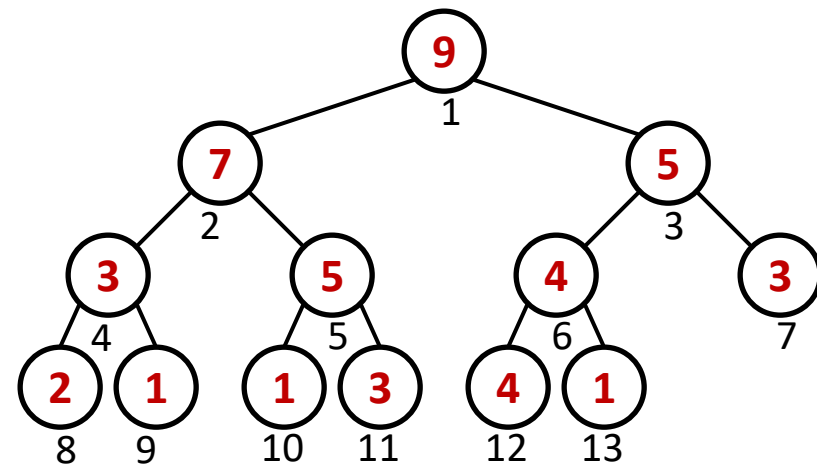
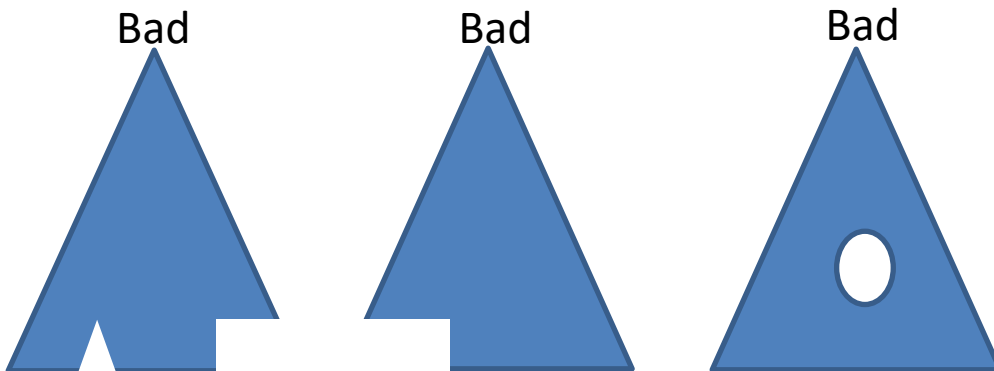
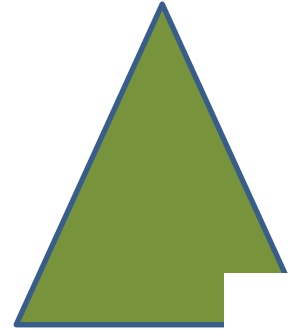
Level	Nodes per level	Sum of nodes from root up to this level
0	$2^0$ (=1)	$2^1 - 1$ (=1)
1	$2^1$ (=2)	$2^2 - 1$ (=3)
2	$2^2$ (=4)	$2^3 - 1$ (=7)
...	...	
$i$	$2^i$	$2^{i+1} - 1$
...	...	
$h-1$	$2^{h-1}$	$2^h - 1$
$h$	$2^h$	$2^{h+1} - 1$



# Nearly Complete Tree

- All levels are full, except possibly for the last level.
- At the last level:
  - Nodes are on the left.
  - Empty positions are on the right.
- There is “no hole”

Good



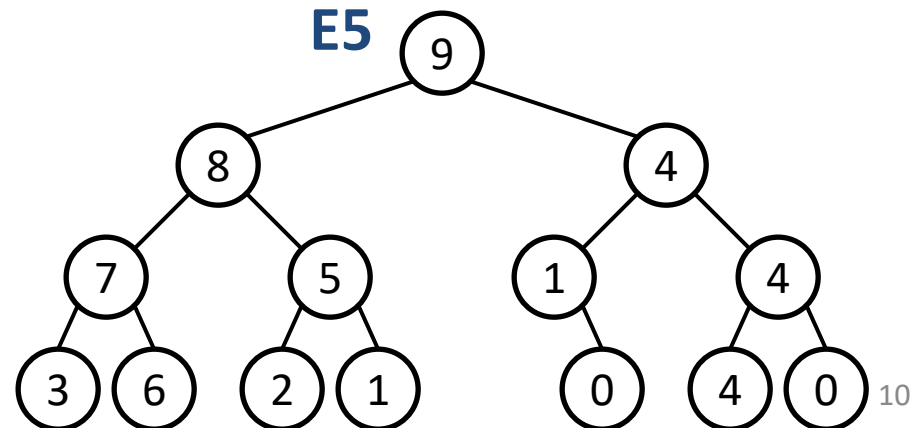
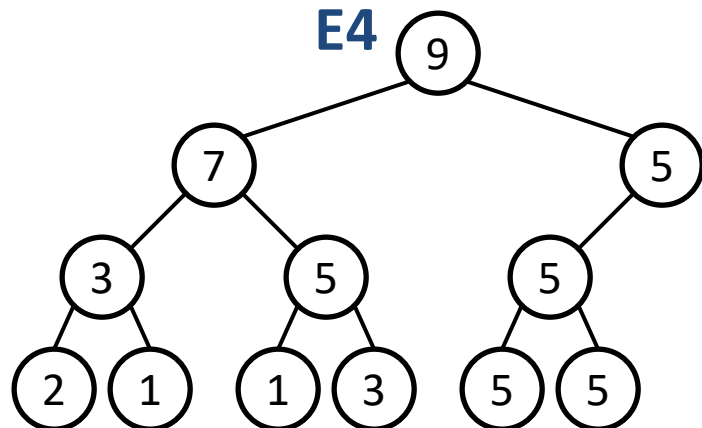
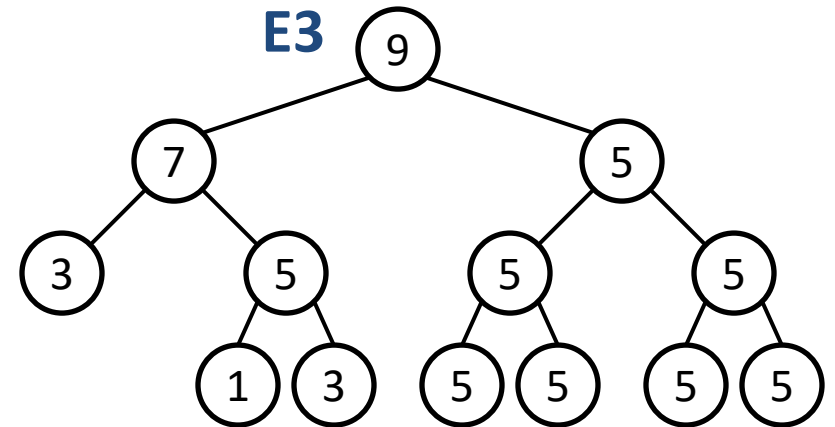
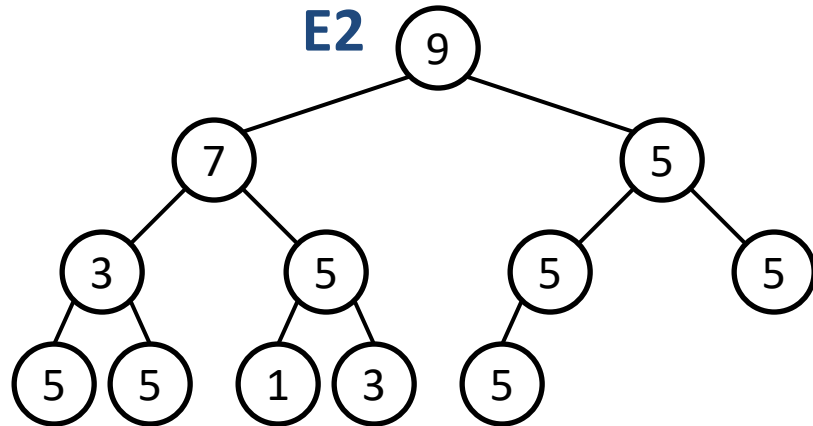
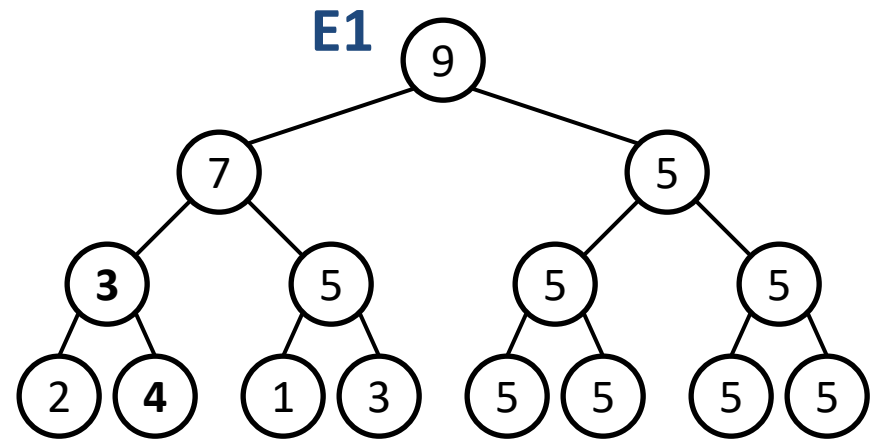


# Worksheet

- Self study: Give examples of trees that are:
  - Complete
  - Full but not nearly complete
  - Nearly complete but not full
  - Neither full nor nearly complete

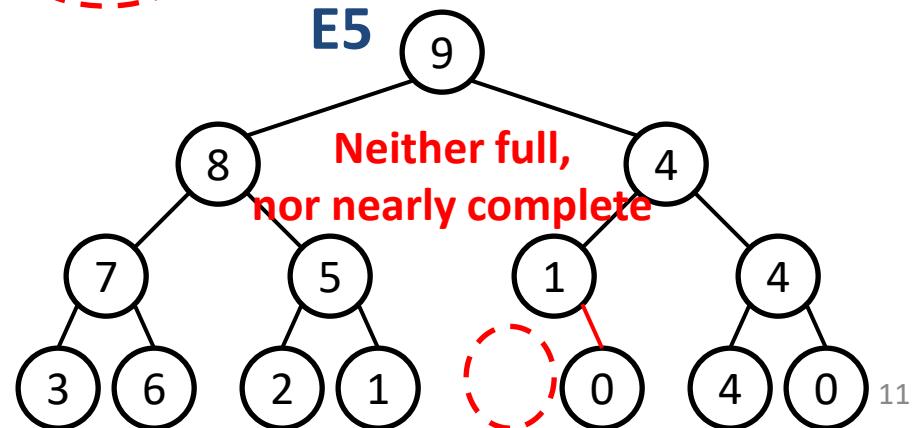
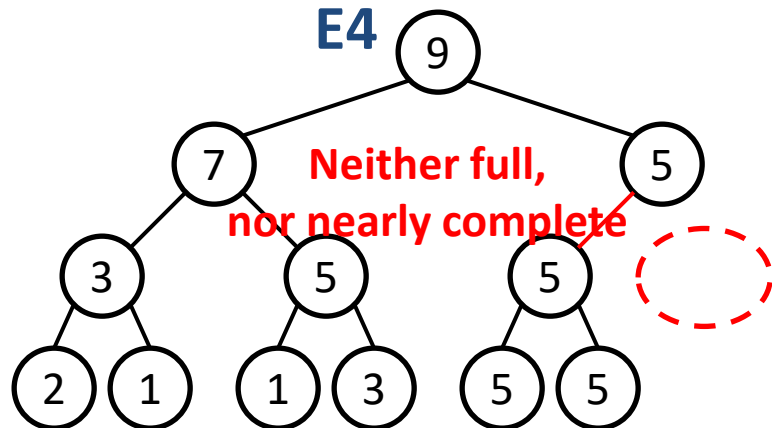
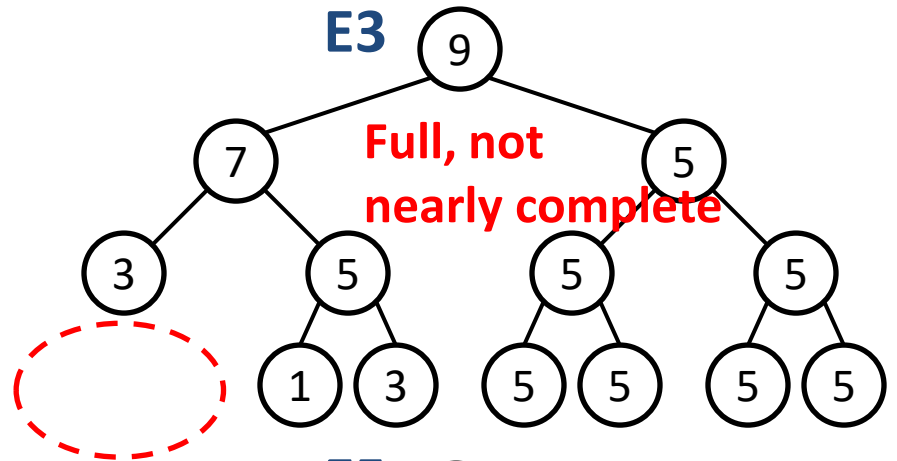
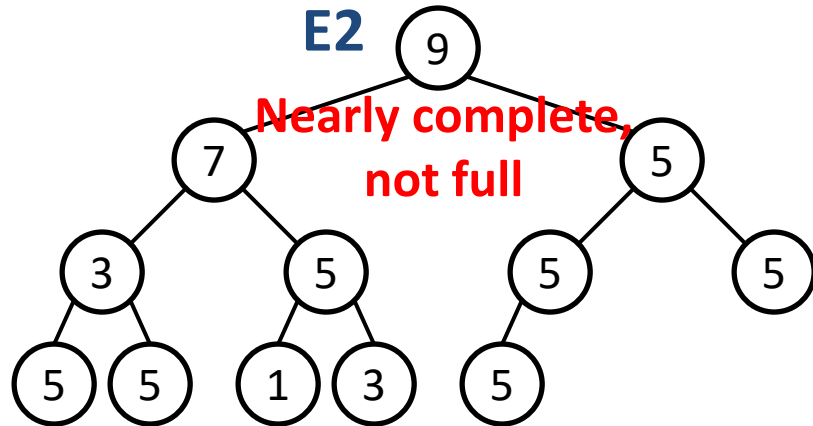
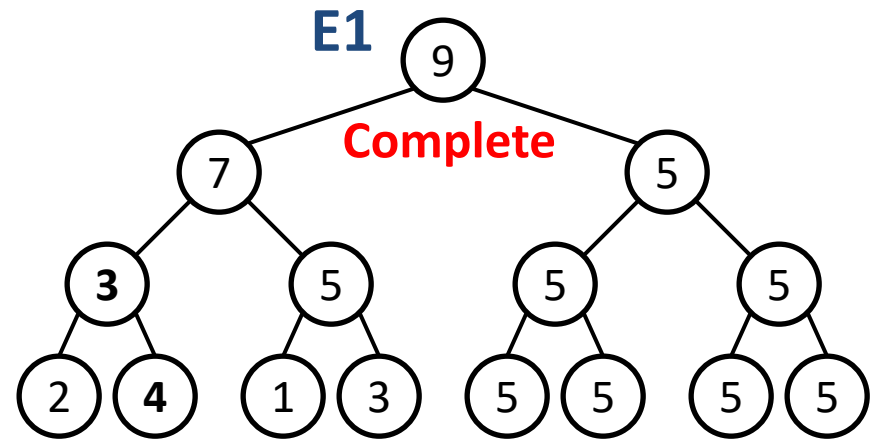
# Worksheet

For each tree, say if it is full, nearly complete or complete.



# Answers

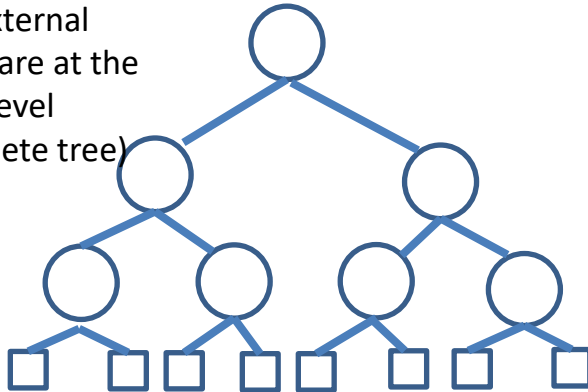
For each tree, say if it is perfect, nearly complete or complete.



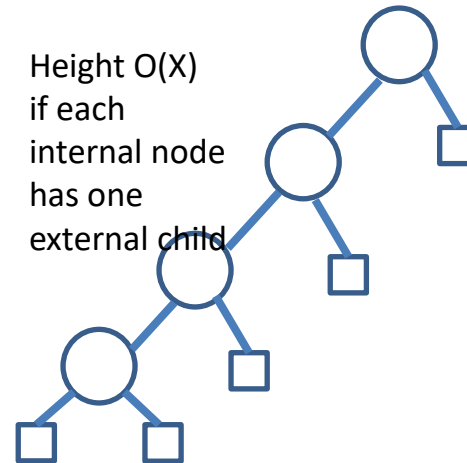
# Properties of Full Trees

- **Full** binary tree : every node has exactly 0 or 2 children. No nodes with only 1 child.
- A **full** binary tree with **X internal nodes** has:
  - $X+1$  external nodes.
  - $2X$  edges (links).
  - $N = 2X+1$  (total number of nodes)
  - height at least  $\lg X$  and at most  $X$ :

Height  $O(\lg X)$   
if all external  
nodes are at the  
same level  
(complete tree)



Height  $O(X)$   
if each  
internal node  
has one  
external child



# Proof

Prove that a full tree with  $P$  internal nodes has  $P+1$  external leaves.

- Full tree property:
- What proof style will you use?

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