

Representation of graph edges, E , : M = matrix, LL = Linked List

Algorithm		Time	Space	Application	Other
BFS	LL	$O(V+E)$	$\Theta(V)$ (color, d, p array, Queue max size)	Flight - fewest connections	Undirected (Ok directed)
	M	$O(V^2)$			
DFS	LL	$O(V+E)$	$\Theta(V)$ (color, st, finish, p arrays rec stack max size)	Graph traversal or search Detect cycles Strongly Connected Components Topological sort	Directed (ok Undirected) Recursion
	M	$O(V^2)$			
Topological sort		Same as DFS	Same as DFS	Order of items in a production line Order of finishing dependent tasks	Directed only No cycles allowed
(Strongly) Connected components		Same as DFS	$\Theta(\text{DFS}) + \Theta(G)$ (needs G^T)	Groups with 2-way communication between any pair within the group	Directed (for undirected - BFS+ restart also works)
MST	LL	$O(V+E \lg V)$ *	$\Theta(V)$	Network layout (e.g. cables for electrical network)	Undirected Greedy, Optimal If priority queue is based on edges, space is $\Theta(E)$
Prim	M	$O(V^2)$	$\Theta(V)$ our method		
SPST (Dijkstra)	LL	$O(V+E \lg V)$ * Same as PRIM	$\Theta(V)$ Same as PRIM	Flight - cheapest cost Driving directions (time, or distance)	Directed Greedy, Optimal
	M	$O(V^2)$			
All pairs SPST		$\Theta(V) * \Theta(\text{SPST})$	$\Theta(V^2)$ (row $i = \text{SPST}(G,i)$)	Same as SPST	Same as SPST. Do SPST(G,v) for every vertex, v

* : if connected: $\Theta(E \lg E) = \Theta(E \lg V)$

** : Assume the graph is connected $\Rightarrow E \geq (V-1)$