

CSE 3302 Notes 10: Implementing Subprograms

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10.1. SEMANTICS OF CALLS AND RETURNS

Pascal-S (and PL/0) reflect typical implementations (without optimization in a one-pass compiler)

10.2. IMPLEMENTING “SIMPLE” SUBPROGRAMS

Historical FORTRAN and COBOL

Simple =

- no recursion
- static allocation
- minimal call-by-value
- call-by-reference for anything non-trivial
- If program starts, it will have enough memory

Sebesta, figure 10.2 is *not* a stack - one activation record per subprogram

Layout of activation record (a.r., stack frame), register conventions mandated by vendor

Who saves registers? - caller or callee?

10.3. IMPLEMENTING SUBPROGRAMS WITH STACK-DYNAMIC LOCAL VARIABLES

PL/0

Calling Sequence:

Caller

```
cal l,a: call procedure a (absolute address) at level l
```

```
cal: begin {generate new block mark}
      s[t + 1] := base(l);      { static link for nested procs }
      s[t + 2] := b;           { dynamic link to caller proc }
      s[t + 3] := p;           { return address }
      b := t + 1;              { new base of stack frame }
      p := a;                  { new program counter }
end;
```

Called

jmp over any instructions for any nested procedures

```
int 0,a : increment t-register by a { a includes 3 slots above }
      int: t := t + a; { allocate stack frame including sl, dl, ra }
```

Return Sequence

Called

```
opr 0,a : execute operation a (=0 here)
      opr: case a of {operator}
            0: begin {return}
                  t := b - 1;           { discard stack frame }
                  p := s[t + 3];        { return to caller }
                  b := s[t + 2];        { old base address }
            end;
```

Caller does nothing special

Pascal-S

Interpreter addressing is done using *display* array to avoid cost of base function.

Recursion causes no difficulty as long as subprograms are not nested (e.g. C)

10.4. NESTED SUBPROGRAMS

A nested subprogram may potentially reference local variables for subprograms that contain it. Two difficulties (that also interact):

1. Calls among nested subprograms at same level.
2. Recursion

For a given variable (under static/lexical scoping), the *most recent* invocation of each containing subprograms is the one whose activation record is needed (diagram in notes 05.3).

(The outermost scope is level 0. Nested scopes have increasing level numbers.)

There is exactly one activation record *per level* in the (ascending) *static chain*.

PL/0

Assumption - non-local data is rarely accessed, simple solution is sufficient

Referencing is based on:

Compiler computes level *difference* to put into `lod`, `sto`, and `cal` instructions

`cal` instruction sets saved static link (`s[t + 1]`) to point at next level a.r.

Level difference is used with loop to follow static links:

```
function base(l: integer): integer;
  var b1: integer;
  begin
    b1 := b; {find base l levels down}
    while l > 0 do
      begin
        b1 := s[b1];
        l := l - 1
      end;
    base := b1
  end {base};
```

Pascal-S (section 4 of report)

Makes no assumption regarding locality of references

Instructions include *absolute* level and offset

No static chain traversals. . . since number of scope levels is small, each access references the *display*.

Trivial, but expensive to create a new display for each call (just use the `base` loop).

Instead, maintain one display incrementally:

Updating display for a call is simple. Replace display entry for the level of the *caller*.

Updating display for a return is messier. When the caller is at a deeper level than the called subprogram, then the same number of display entries recovered from the static chain. (This assumes that subsequent calls had occurred.)

(Caches on modern machines lessen the performance advantage of displays)

10.5. BLOCKS

Notes 5, page 3

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PL/0 examples: Lab 4 Spring 2013 - added call-by-value arguments

```

bash-3.2$ cat recurse.pl0 - |plzero.argstack
0 var aout,result2;
1   3   4
1 procedure a(ain);
2   -1
2   var bout;
2   3
2   procedure b(bin);
3     -1
3     var cout;
3     3
3     procedure c(cin);
4       -1
4       begin
5         cout:=cin+8;
9         result2:=ain+bin+cin+8
15      end;
4 int  0   3          c code
5 lod  0  -1 cin
6 lit  0   8
7 opr  0   2
8 sto  1   3 cout
9 lod  2  -1 ain
10 lod  1  -1 bin
11 opr  0   2
12 lod  0  -1 cin
13 opr  0   2
14 lit  0   8
15 opr  0   2
16 sto  3   4 result2
17 opr  0   0
18
18   begin          b code
19     call c(4);
22     bout:=bin+cout
24   end;
18 int  0   4
19 lit  0   4
20 cal  0   4 c
21 int  0  -1
22 lod  0  -1 bin
23 lod  0   3 cout
24 opr  0   2
25 sto  1   3 bout
26 opr  0   0
27
27   begin          a code
28     call b(2);
31     aout:=ain+bout
33   end;
27 int  0   4
28 lit  0   2
29 cal  0  18 b
30 int  0  -1
31 lod  0  -1 ain
32 lod  0   3 bout
33 opr  0   2
34 sto  1   3 aout
35 opr  0   0
36
36 begin
37 call a(1);
40 out:=aout;
42 out:=result2
43 end.
36 int  0   5
37 lit  0   1
38 cal  0  27 a
39 int  0  -1
40 lod  0   3 aout
41 wro  0   0
42 lod  0   4 result2
43 wro  0   0
44 opr  0   0

```

```

start pl/0
t=5 b=1 p=37      initial
s[1]=0  s.l.
s[2]=0  d.l.
s[3]=0  ret adr
s[4]=0  aout
s[5]=0  result2
t=10 b=7 p=28    after call to a
s[1]=0  s.l.
s[2]=0  d.l.
s[3]=0  ret adr
s[4]=0  aout
s[5]=0  result2
s[6]=1  ain
s[7]=1  s.l.
s[8]=1  d.l.
s[9]=39  ret adr
s[10]=0  bout
t=15 b=12 p=19   after call to b
s[1]=0  s.l.
s[2]=0  d.l.
s[3]=0  ret adr
s[4]=0  aout
s[5]=0  result2
s[6]=1  ain
s[7]=1  s.l.
s[8]=1  d.l.
s[9]=39  ret adr
s[10]=0  bout
s[11]=2  bin
s[12]=7  s.l.
s[13]=7  d.l.
s[14]=30  ret adr
s[15]=0  cout
t=19 b=17 p=5    after call to c
s[1]=0  s.l.
s[2]=0  d.l.
s[3]=0  ret adr
s[4]=0  aout
s[5]=0  result2
s[6]=1  ain
s[7]=1  s.l.
s[8]=1  d.l.
s[9]=39  ret adr
s[10]=0  bout
s[11]=2  bin
s[12]=7  s.l.
s[13]=7  d.l.
s[14]=30  ret adr
s[15]=0  cout
s[16]=4  cin
s[17]=12 s.l.
s[18]=12 d.l.
s[19]=21  ret adr
! 15
! 15
end pl/0

```

```

-bash-3.2$ cat rtsExample.pl0 - |
plzero.argstack
0 procedure iloop(i);
2     -1
2 procedure jloop(j);
3     -1
3 procedure kloop(k);
4     -1
4     begin
5         k:=k+1;
9         out:=i+j+k;
15        if k<3 then
18            call kloop(k);
22        if k=3 then
25            call jloop(j)
29        end;
4 int 0 3
5 lod 0 -1 k
6 lit 0 1
7 opr 0 2
8 sto 0 -1 k
9 lod 2 -1 i
10 lod 1 -1 j
11 opr 0 2
12 lod 0 -1 k
13 opr 0 2
14 wro 0 0
15 lod 0 -1 k
16 lit 0 3
17 opr 0 10
18 jpc 0 22
19 lod 0 -1 k
20 cal 1 4 kloop
21 int 0 -1
22 lod 0 -1 k
23 lit 0 3
24 opr 0 8
25 jpc 0 29
26 lod 1 -1 j
27 cal 2 2 jloop
28 int 0 -1
29 opr 0 0
30
30     begin jloop code
31         j:=j+10;
35         if j<40 then
38             call kloop(0);
42         if j=40 then
45             call iloop(i)
49         end;
30 int 0 3
31 lod 0 -1 j
32 lit 0 10
33 opr 0 2
34 sto 0 -1 j
35 lod 0 -1 j
36 lit 0 40
37 opr 0 10
38 jpc 0 42
39 lit 0 0
40 cal 0 4 kloop
41 int 0 -1
42 lod 0 -1 j
43 lit 0 40
44 opr 0 8
45 jpc 0 49
46 lod 1 -1 i
47 cal 2 1 iloop
48 int 0 -1
49 opr 0 0
50
50     begin iloop code
51         i:=i+100;
55         if i<400 then
58             call jloop(0)
62         end;
50 int 0 3
51 lod 0 -1 i
52 lit 0 100
53 opr 0 2
54 sto 0 -1 i
55 lod 0 -1 i
56 lit 0 400
57 opr 0 10
58 jpc 0 62
59 lit 0 0
60 cal 0 30 jloop
61 int 0 -1
62 opr 0 0
63
63 begin
64     call iloop(0)
67 end.
63 int 0 3
64 lit 0 0
65 cal 0 50 iloop
66 int 0 -1
67 opr 0 0
start pl/0
...
t=167 b=165 p=5 (beg kloop)
s[1]=0 s.l.
s[2]=0 d.l.
s[3]=0
s[4]=100
s[5]=1
s[6]=1 d.l.
s[7]=66
s[8]=10
s[9]=5
s[10]=5 d.l.
s[11]=61
s[12]=1
s[13]=9
s[14]=9 d.l.
s[15]=41
s[16]=2
s[17]=9
s[18]=13 d.l.
s[19]=21
s[20]=3
s[21]=9
s[22]=17 d.l.
s[23]=21
s[24]=20
s[25]=5
s[26]=21 d.l.
s[27]=28
s[28]=1
s[29]=25
s[30]=25 d.l.
s[31]=41
s[32]=2
s[33]=25
s[34]=29 d.l.
s[35]=21
s[36]=3
s[37]=25
s[38]=33 d.l.
s[39]=21
s[40]=30
s[41]=5
s[42]=37 d.l.
s[43]=28
s[44]=1
s[45]=41
s[46]=41 d.l.
s[47]=41
s[48]=2
s[49]=41
s[50]=45 d.l.
s[51]=21
s[52]=3
s[53]=41
s[54]=49 d.l.
s[55]=21
s[56]=40
s[57]=5
s[58]=53 d.l.
s[59]=28
s[60]=200
s[61]=1
s[62]=57 d.l.
s[63]=48
s[64]=10
s[65]=61
s[66]=61 d.l.
s[67]=61
s[68]=1
s[69]=65
s[70]=65 d.l.
s[71]=41
s[72]=2
s[73]=65
s[74]=69 d.l.
s[75]=21
s[76]=3
s[77]=65
s[78]=73 d.l.
s[79]=21
s[80]=20
s[81]=61
s[82]=77 d.l.
s[83]=28
s[84]=1
s[85]=81
s[86]=81 d.l.
s[87]=41
s[88]=2
s[89]=81
s[90]=85 d.l.
s[91]=21
s[92]=3
s[93]=81
s[94]=89 d.l.
s[95]=21
s[96]=30
s[97]=61
s[98]=93 d.l.
s[99]=28
s[100]=1
s[101]=97
s[102]=97 d.l.
s[103]=41
s[104]=2
s[105]=97
s[106]=101 d.l.
s[107]=21
s[108]=3
s[109]=97
s[110]=105 d.l.
s[111]=21
s[112]=40
s[113]=61
s[114]=109 d.l.
s[115]=28
s[116]=300 i
s[117]=1 s.l.
s[118]=113 d.l.
s[119]=48
s[120]=10
s[121]=117
s[122]=117 d.l.
s[123]=61
s[124]=1
s[125]=121
s[126]=121 d.l.
s[127]=41
s[128]=2
s[129]=121
s[130]=125 d.l.

```

```
s[131]=21
s[132]=3
s[133]=121
s[134]=129 d.l.
s[135]=21
s[136]=20
s[137]=117
s[138]=133 d.l.
s[139]=28
s[140]=1
s[141]=137
s[142]=137 d.l.
s[143]=41
s[144]=2
s[145]=137
s[146]=141 d.l.
s[147]=21
s[148]=3
s[149]=137
s[150]=145 d.l.
s[151]=21
s[152]=30 j
s[153]=117 s.l.
s[154]=149 d.l.
s[155]=28 ret adr (exit kloop)
s[156]=1
s[157]=153
s[158]=153 d.l.
s[159]=41
s[160]=2
s[161]=153
s[162]=157 d.l.
s[163]=21
s[164]=2 k
s[165]=153 s.l.
s[166]=161 d.l.
s[167]=21 ret adr (if k=3 then)
! 333
```