## CSE 3318 Notes 3: Summations

(Last updated 1/5/24 10:26 AM)

CLRS, appendix A

3.A. GEOMETRIC SERIES (review)

 $\sum_{k=0}^{t} x^{k} = \frac{x^{t+1} - 1}{x - 1} \quad \text{when } x \neq 1 \quad \text{[Not hard to verify by math induction]}$ 

$$\sum_{k=0}^{t} x^{k} \le \sum_{k=0}^{\infty} x^{k} = \lim_{k \to \infty} \frac{x^{k} - 1}{x - 1} = \frac{1}{1 - x} \quad \text{when } 0 < x < 1$$

**3.B.** HARMONIC SERIES

$$\ln n \le H_n = \sum_{k=1}^n \frac{1}{k} \le \ln n + .577... \le \ln n + 1$$

3.C. APPROXIMATION BY INTEGRALS (p. 1150-1151)

For a monotonically increasing  $(x \le y \Rightarrow f(x) \le f(y))$  function:

$$\int_{m-1}^{n} f(x)dx \leq \sum_{k=m}^{n} f(k) \leq \int_{m}^{n+1} f(x)dx$$

Since:

$$\begin{split} & \underset{k-1}{\overset{k}{\int}} f(x) dx \leq f(k) \leq \underset{k}{\overset{k+1}{\int}} f(x) dx \\ \end{split}$$

in this situation.

## 3.D. BOUNDING SUMMATIONS USING MATH INDUCTION AND INEQUALITIES

[Techniques are especially important for recurrences in Notes 04]

Show 
$$\sum_{i=1}^{n} i^2 = \Theta(n^3)$$
 [Trivial to show using integration or  $\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$ .]

a. Show  $O(n^3)$ 

(i) 
$$\sum_{i=1}^{n=1} i^2 = 1 \le cn^3$$
 using any constant  $c \ge 1$ 

(ii) Suppose this holds for *n*:

$$\sum_{i=1}^{n} i^2 \le cn^3$$

Now go on to n + 1 and show that the bound *still holds* 

$$\sum_{i=1}^{n+1} i^{2} = \sum_{i=1}^{n} i^{2} + (n+1)^{2}$$
$$= \sum_{i=1}^{n} i^{2} + n^{2} + 2n + 1$$
$$\leq cn^{3} + n^{2} + 2n + 1$$
$$= ???$$
$$\leq c(n+1)^{3}$$

The bridging step (???) separates the bounding term ( $c(n+1)^3$ ) from everything else (x):

$$c(n+1)^{3} + x = cn^{3} + n^{2} + 2n + 1$$
  

$$x = cn^{3} + n^{2} + 2n + 1 - cn^{3} - 3cn^{2} - 3cn - c = (1 - 3c)n^{2} + (2 - 3c)n + 1 - c$$
  
So ??? is now  $c(n+1)^{3} + [(1 - 3c)n^{2} + (2 - 3c)n + 1 - c]$ 

Can drop [ . . . ] (through  $\leq$ ) if it cannot become positive. Happens if  $c \geq 1$ 

b. Show  $\Omega(n^3)$ 

(i) 
$$\sum_{i=1}^{n=1} i^2 = 1 \ge cn^3$$
 using any constant  $0 < c \le 1$ 

(ii) Suppose this holds for *n*:

$$\sum_{i=1}^{n} i^2 \ge cn^3$$

Now go on to n + 1 and show that the bound *still holds* 

$$\sum_{i=1}^{n+1} i^{2} = \sum_{i=1}^{n} i^{2} + (n+1)^{2}$$
$$= \sum_{i=1}^{n} i^{2} + n^{2} + 2n + 1$$
$$\ge cn^{3} + n^{2} + 2n + 1$$
$$= ???$$
$$\ge c(n+1)^{3}$$

The bridging step (???) involves the same algebra as before.

Can drop [...] (through  $\geq$ ) if it cannot become negative. Happens if  $0 < c \le 1/3$ 

Suppose we attempt to show  $\sum_{i=1}^{n} i^2 = \Theta(n^2)$ 

a. Show  $O(n^2)$ 

(i) 
$$\sum_{i=1}^{n=1} i^2 = 1 \le cn^2$$
 using any constant  $c \ge 1$ 

(ii) Suppose this holds for *n*:

$$\sum_{i=1}^{n} i^2 \le cn^2$$

Now attempt to go on to n + 1.

$$\sum_{i=1}^{n+1} i^2 = \sum_{i=1}^n i^2 + (n+1)^2$$
$$= \sum_{i=1}^n i^2 + n^2 + 2n + 1$$
$$\leq cn^2 + n^2 + 2n + 1$$
$$= ???$$
$$\leq c(n+1)^2$$

The bridging step (???) separates the bounding term  $(c(n+1)^2)$  from everything else (x):

$$c(n+1)^{2} + x = cn^{2} + n^{2} + 2n + 1$$
  

$$x = cn^{2} + n^{2} + 2n + 1 - cn^{2} - 2cn - c = n^{2} + (2 - 2c)n + 1 - c$$
  
So ??? is now  $c(n+1)^{2} + \left[n^{2} + (2 - 2c)n + 1 - c\right]$ 

Can drop [ . . . ] (through  $\leq$ ) if it cannot become positive. *Fails as n grows.* 

b. Can still show  $\Omega(n^2)$ 

(i) 
$$\sum_{i=1}^{n=1} i^2 = 1 \ge cn^2$$
 using any constant  $0 < c \le 1$ 

(ii) Suppose this holds for *n*:

$$\sum_{i=1}^{n} i^2 \ge cn^2$$

Now go on to n + 1.

$$\sum_{i=1}^{n+1} i^2 = \sum_{i=1}^n i^2 + (n+1)^2$$
$$= \sum_{i=1}^n i^2 + n^2 + 2n + 1$$
$$\ge cn^2 + n^2 + 2n + 1$$
$$= ???$$
$$\ge c(n+1)^2$$

The bridging step separates the bounding term  $(c(n+1)^2)$  from everything else (*x*):

$$c(n+1)^{2} + x = cn^{2} + n^{2} + 2n + 1$$
  

$$x = cn^{2} + n^{2} + 2n + 1 - cn^{2} - 2cn - c = n^{2} + (2 - 2c)n + 1 - c$$
  
So ??? is now  $c(n+1)^{2} + [n^{2} + (2 - 2c)n + 1 - c]$ 

Can drop [ . . . ] (through  $\geq$ ) if it cannot become negative.

Happens if  $0 < c \le 1$  (or for "sufficiently large" *n*).